

COMP3012/G53CMP: Lecture 6

Contextual Analysis: Scope II

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Identification for LTXL

We are now going to study a concrete Haskell implementation of identification for **LTXL**:

Less Trivial expression Language

- LTXL \approx TXL + typed definitions + if-expression + new operators
- Slides only show highlights: complete code available on-line.

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LTXL CFG (3)

```
PrimaryExp → LitInt
            | Ident
            | \ PrimaryExp
            | - PrimaryExp
            | if Exp then Exp else Exp
            | ( Exp )
            | let Defs in Exp
```

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This Lecture

An Illustrative Identification Algorithm in Haskell

- LTXL Syntax and Semantics, particularly scope rules.
- Abstract syntax representation
- Environment/Symbol Table representation and operations.
- The Identification Algorithm

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LTXL CFG (1)

$LTXLProgram \rightarrow Exp$

```
Exp → Exp || Exp | Exp && Exp
    | Exp < Exp | Exp == Exp | Exp > Exp
    | Exp + Exp | Exp - Exp
    | Exp * Exp | Exp / Exp
    | PrimaryExp
```

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LTXL CFG (4)

```
Defs → Def ; Defs
      | Def
Def → Type Ident = Exp
Type → int
      | bool
```

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Recap: Identification

Identification is the task of relating each applied identifier occurrence to its declaration or definition:

```
public class C {
    int x, n;
    void set(int n) { x = n; }
}
```

In the body of `set`, the one applied occurrence of

- `x` refers to the **instance variable** `x`
- `n` refers to the **argument** `n`.

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LTXL CFG (2)

Operator precedence and associativity is used to disambiguate. In increasing order of precedence:

1. `||`
2. `&&`
3. `<, ==, >`
4. `+, -`
5. `*, /`

All left associative.

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LTXL Example 1

```
let
    int a = 10;
    bool b = a < 2
in let
    int c = a * 10;
    bool a = a == 42;
    int d = if a then 1 else 2
in
    if a && b then c else 42
```

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LTXL Scope Rules

1. The scope of a variable is all subsequent definitions and the body of the `let`-expression in which the definition of the variable occurs. A variable is **not** in scope in the RHS of its own definition.
2. A definition of a variable hides, for the extent of its scope, any definition of a variable with the same name from an outer `let`-expression.
3. At most one definition may be given for a variable in the list of definitions of a `let`-expression.

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LTXL Example 1 (again)

Which scope rules are used where?

```
let
  int a = 10;
  bool b = a < 2
in let
  int c = a * 10;
  bool a = a == 42;
  int d = if a then 1 else 2
in
  if (a && b) then c else 42
```

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LTXL Example 2

What about this LTXL example?

```
let
  int a = 1;
  int b = c * 2;
  bool a = a < 1
in
  a + b
```

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LTXL AST (1)

The following Haskell data types are used to represent LTXL programs.

```
type Id = String

data Type = IntType
          | BoolType
          | UnknownType
```

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LTXL AST (2)

```
data UnOp = Not | Neg

data BinOp = Or
           | And
           | Less
           | Equal
           | Greater
           | Plus
           | Minus
           | Times
           | Divide
```

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LTXL AST (3)

`Exp` is a **parameterized** type. The **type parameter** `a` allows variables to be **annotated** with an attribute of type `a`. This facility is used by the identification function.

```
data Exp a
  = LitInt Int
  | Var Id a
  | UnOpApp UnOp (Exp a)
  | BinOpApp BinOp (Exp a) (Exp a)
  | If (Exp a) (Exp a) (Exp a)
  | Let [(Id, Type, Exp a)] (Exp a)
```

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LTXL AST (4)

Example: The LTXL program

```
let int x = 7 in x + 35
```

would be represented like this **before** identification (type `Exp ()`):

```
Let [("x", IntType, LitInt 7)]
  (BinOpApp Plus
   (Var "x" ())
   (LitInt 35))
```

(**After** identification, type will be `Exp Attr`.)

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LTXL Environment (1)

- An **association list** is used to represent the environment/symbol table to keep things simple.
- By **prepending** new declarations to the list, and searching from the beginning, we will always find an identifier in the closest containing scope. For example:

```
lookup "x" [("x", a1), ("y", a2), ("x", a3)]
⇒ a1
```
- No need for a "close scope" operation. We are in a pure functional setting ⇒ persistent data.

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LTXL Environment (2)

The environment associates identifiers with **variable attributes**. Our attributes are the **scope level** and the **declared type**.

```
type Attr = (Int, Type)
```

The environment is just an association list:

```
type Env = [(Id, Attr)]
```

Note: our environment does **not** store variable **definitions**.

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LTXL Environment (3)

Example:

```
let
  int a = 10;           (1)
  int b = a + 42
in let
  bool a = b < 20      (2)
in
  if a then b else 13
```

Env. after (1): [("a", (1, IntType))]
Env. after (2): [("a", (2, BoolType)),
("b", (1, IntType)), ("a", (1, IntType))]

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LTXL Environment (5)

```
enterVar i l t env
| not (isDefined i l env) -- decl. prepended
= Left (i, (l,t)) : env
| otherwise
= Right (i ++ " already defined.")
where
  isDefined i l [] = False
  isDefined i l ((i', (l',_)) : env)
  | l < l' = error "Should not happen!"
  | l > l' = False
  | i == i' = True
  | otherwise = isDefined i l env
```

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LTXL Environment (8)

Let

```
env = [ ("x", (2, BoolType)),
        ("y", (2, IntType)),
        ("x", (1, IntType)) ]
```

Then:

```
lookupVar "y" env
⇒ Left (2, IntType)
lookupVar "x" env
⇒ Left (2, BoolType)
lookupVar "z" env
⇒ Right "z not defined."
```

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LTXL Environment (4)

enterVar inserts a variable at the given scope level and of the given type into an environment.

- Check that no variable with same name has been defined at the same scope level.
- If not, the new variable is entered, and the **resulting environment** is returned.
- Otherwise an **error message** is returned.

```
enterVar :: Id -> Int -> Type -> Env
          -> Either Env ErrorMsg
```

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LTXL Environment (6)

Let

```
env = [ ("y", (2, IntType)),
        ("x", (1, IntType)) ]
```

Then:

```
enterVar "x" 2 BoolType env
⇒ Left [ ("x", (2, BoolType)),
         ("y", (2, IntType)),
         ("x", (1, IntType)) ]
```

```
enterVar "y" 2 BoolType env
⇒ Right "y already defined."
```

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LTXL Identification (1)

Goals of LTXL identification phase:

- Annotate each applied identifier occurrence with attributes of the corresponding variable declaration. I.e., map unannotated AST **Exp ()** to annotated AST **Exp Attr**.
- Report conflicting variable definitions and undefined variables.

```
identification :: Exp () -> (Exp Attr, [ErrorMsg])
```

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Aside: The Haskell Type Either

The standard Haskell type `Either` comes in handy when one needs to represent a value that has one of two possible types:

```
data Either a b = Left a | Right b
```

A typical example is when a function needs to return one of two kinds of results:

```
foo :: Int -> Either Bool String
foo x | x < 100 = Left (x < 0)
      | otherwise = Right "Too big"
```

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LTXL Environment (7)

lookupVar looks up a variable in an environment.

- Returns **variable attributes** if found.
- Returns an **error message** otherwise.

```
lookupVar :: Id -> Env
          -> Either Attr ErrorMsg
lookupVar i [] = Right (i ++ " not defined.")
lookupVar i ((i', a) : env)
  | i == i' = Left a
  | otherwise = lookupVar i env
```

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LTXL Identification (2)

Example: Before Identification

```
Let [("x", IntType, LitInt 7)]
  (BinOpApp Plus
   (Var "x" ())
   (LitInt 35))
```

After identification:

```
Let [("x", IntType, LitInt 7)]
  (BinOpApp Plus
   (Var "x" (1, IntType))
   (LitInt 35))
```

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LTXL Identification (3)

Main identification function:

```
identification :: Exp ()
               -> (Exp Attr, [ErrorMsg])
identification e = identAux 0 emptyEnv e
```

Type signature for auxiliary identification function:

```
identAux :: Int -> Env -> Exp ()
         -> (Exp Attr, [ErrorMsg])
```

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LTXL Identification (6)

Reminder: LTXL scope rules

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2. A definition of a variable hides, for the extent of its scope, any definition of a variable with the same name from an outer `let`-expression.
3. At most one definition may be given for a variable in the list of definitions of a `let`-expression.

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Efficient Symbol Table Implementation

Lists don't make for very efficient symbol tables. Insertion (at head) is fast, $O(1)$, but lookup is $O(n)$, where n is the number of symbols.

Some more efficient options:

- Balanced trees:
 - Insertion and lookup are both $O(\log n)$.
 - One way of handling nested scopes would be a stack of trees.

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LTXL Identification (4)

Variable case:

```
identAux l env (Var i _) =
  case lookupVar i env of
  Left a  -> (Var i a, [])
  Right m -> (Var i (0, UnknownType), [m])
```

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LTXL Identification (7)

Block of definitions (`let`):

```
identAux l env (Let ds e) =
  (Let ds' e', ms1 ++ ms2)
  where
    l' = l + 1
    (ds', env', ms1) = identDefs l' env ds
    (e', ms2)        = identAux l' env' e
```

Note that `identDefs` returns an **updated environment** to be **used** when checking the **body** of the `let` (rule 1).

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Efficient Symbol Table Implementation

• Hash tables:

- Insertion and lookup are both $O(1)$ as long as the ratio between the number of symbols and the hash table size is kept below a small constant factor.
- Algorithms such as **linear hashing** allows the table to grow and shrink gracefully, guaranteeing near optimal performance.

See e.g. Aho, Sethi, Ullman (1986) for further details.

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LTXL Identification (5)

Binary operator application (typical recursive case):

```
identAux l env (BinOpApp op e1 e2) =
  (BinOpApp op e1' e2', ms1 ++ ms2)
  where
    (e1', ms1) = identAux l env e1
    (e2', ms2) = identAux l env e2
```

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LTXL Identification (8)

```
identDefs l env [] = ([], env, [])
identDefs l env ((i,t,e) : ds) =
  ((i,t,e') : ds', env'', ms1++ms2++ms3)
  where
    (e', ms1) = identAux l env e
    (env'', ms2) = impl/checks rules 2 & 3
    case enterVar i l t env of
      Left env' -> (env', [])
      Right m   -> (env, [m])
    (ds', env'', ms3) = i in scope (rule 1)
    identDefs l env' ds
```

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