

G54FOP: Lecture 6

Operational Semantics III: State

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Introducing State: Key Ideas

- μ : Store (state, memory); can be read/updated.
- Evaluation relation relates terms **and stores** to new terms and stores:

$$t \mid \mu \longrightarrow t' \mid \mu'$$

- $l \in \mathcal{L}$: Uninterpreted set of locations (addresses) with equality.
- $\mu \in \mathcal{L} \rightarrow v$: Store: map from location to value.
- $\mu(l)$: Lookup value at location l .
- $[l \mapsto v]\mu$: Update; like μ except $([l \mapsto v]\mu)(l) = v$

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Homework Lecture 6 (1)

1. Consider the Small Imperative Language. Add a loop construct:

$$t \rightarrow \dots$$

$\mathbf{while} \ t \ \mathbf{do} \ t$	<i>while loop</i>
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Provide evaluation rule(s) for this construct, assuming the usual semantics of a while loop: repetition of loop body **zero** or more times as long as loop condition is true.

Hint: Make use of what you have!

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Small Expression Language: Terms

t	\rightarrow \mathbf{true} \mathbf{false} $\mathbf{if} \ t \ \mathbf{then} \ t \ \mathbf{else} \ t$ 0 $\mathbf{succ} \ t$ $\mathbf{pred} \ t$ $\mathbf{iszero} \ t$	<i>terms:</i> <i>constant true</i> <i>constant false</i> <i>conditional</i> <i>constant zero</i> <i>successor</i> <i>predecessor</i> <i>zero test</i>
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Small Imperative Language (1)

New terms; extends the terms of Small Expression Language:

t	\rightarrow \dots \mathbf{unit} $\mathbf{new} \ t$ $\mathbf{!} \ t$ $t \ \mathbf{:=} \ t$ l $t \ ; \ t$	<i>terms:</i> <i>constant unit</i> <i>allocation</i> <i>dereferencing</i> <i>assignment</i> <i>store location</i> <i>sequencing</i>
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Homework Lecture 6 (2)

2. As mentioned, our language is still an expression language where expressions may have side effects. Design a new language (syntax and op. sem.) by separating the terms into
 - expressions: do not have side effects
 - commands: have side effects
 and making any other changes you see fit. Don't worry too much if the resulting language isn't "useful" (the Small Imperative Language isn't really). Can the evaluation rules for expressions somehow be simplified by exploiting that expressions do not have side effects?

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Small Expression Language: Values

v	\rightarrow \mathbf{true} \mathbf{false} nv nv 0 $\mathbf{succ} \ nv$	<i>values:</i> <i>constant true</i> <i>constant false</i> <i>numeric value</i> <i>numeric values:</i> <i>zero value</i> <i>successor value</i>
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Small Imperative Language (2)

New values; extends the values of Small Expression Language:

v	\rightarrow \dots \mathbf{unit} l	<i>values:</i> <i>unit value</i> <i>store location</i>
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Note: Still an expression language in that every term is an expression that evaluates to a value, even if some expressions have side effects. No separate category of commands.

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