G54FOP: Lecture 16 Denotational Semantics and Domain Theory II

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This Lecture

- Denotational semantics for small imperative language.
- Introduction to semantics of loops and recursion.

Imperative Language (1)

Syntax of expressions:

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expressions:		\rightarrow	0
variable	${\mathcal X}$		
constant natural number, $\mathbb N$	n		
constant true	true		
constant false	false		
logical negation	not e		
logical conjunction	e 🍇 e		

Imperative Language (2)



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Imperative Language (3)

Syntax of commands:

C

 \rightarrow commands:| skipno operation| x := eassignment| c ; csequence| if e then c else cconditionalwhile e do citeration

Semantics of Expressions (1)

We take the *semantic domain* to be \mathbb{N} for simplicity. We need a way to give meaning to *variables*. A *store* maps a variable name to its value:

$$\begin{array}{rcl} \Sigma &=& x \to \mathbb{N} \\ \sigma & : & \Sigma \end{array}$$

Semantics of Expressions (2)

We then need two **semantic functions**, one for expressions (have no side effects in this language), one for commands.

Starting with the one for expressions:

 $\mathbb{E}\llbracket \cdot \rrbracket : e \to (\Sigma \to \mathbb{N})$ (Note: $e \to (\Sigma \to \mathbb{N}) = e \to \Sigma \to \mathbb{N}$ etc.)

(Definition on whiteboard)

Semantics of Commands

A command is executed for its *effects*: given a state, executing a command results in a new state. A command is a *state transformer*.

In our case, the state comprises only the store:

 $\Sigma = x \to \mathbb{N}$

Thus, type of state transformer: $\Sigma \to \Sigma$. Semantic function for commands: $C[\cdot] : c \to (\Sigma \to \Sigma)$ [Not correct yet!] (Definition on whiteboard)