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The Arpeggigon: A Functional Reactive Musical Automaton

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The Arpeggigon (1)

- Software realisation of the reacTogon:



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- Interactive cellular automaton:
 - Configuration
 - Performance parameters

The Arpeggigon (2)

- Implemented in Haskell using:
 - The Functional Reactive Programming (FRP) system Yampa
 - Reactive Values and Relations (RVR)

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Code: <https://gitlab.com/chupin/arpeggigon>

Video:

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Before you get too excited: ***Work in progress!***

Motivation

Exploring FRP and RVR as an (essentially) declarative way for developing full-fledged musical applications:

- FRP aligns with declarative and temporal (discrete and continuous) nature of music
- RVR allows declarative-style interfacing with external components

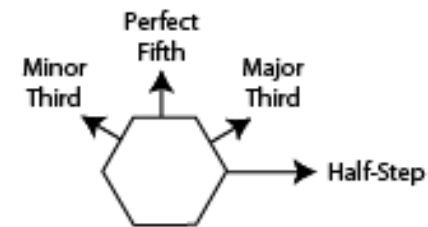
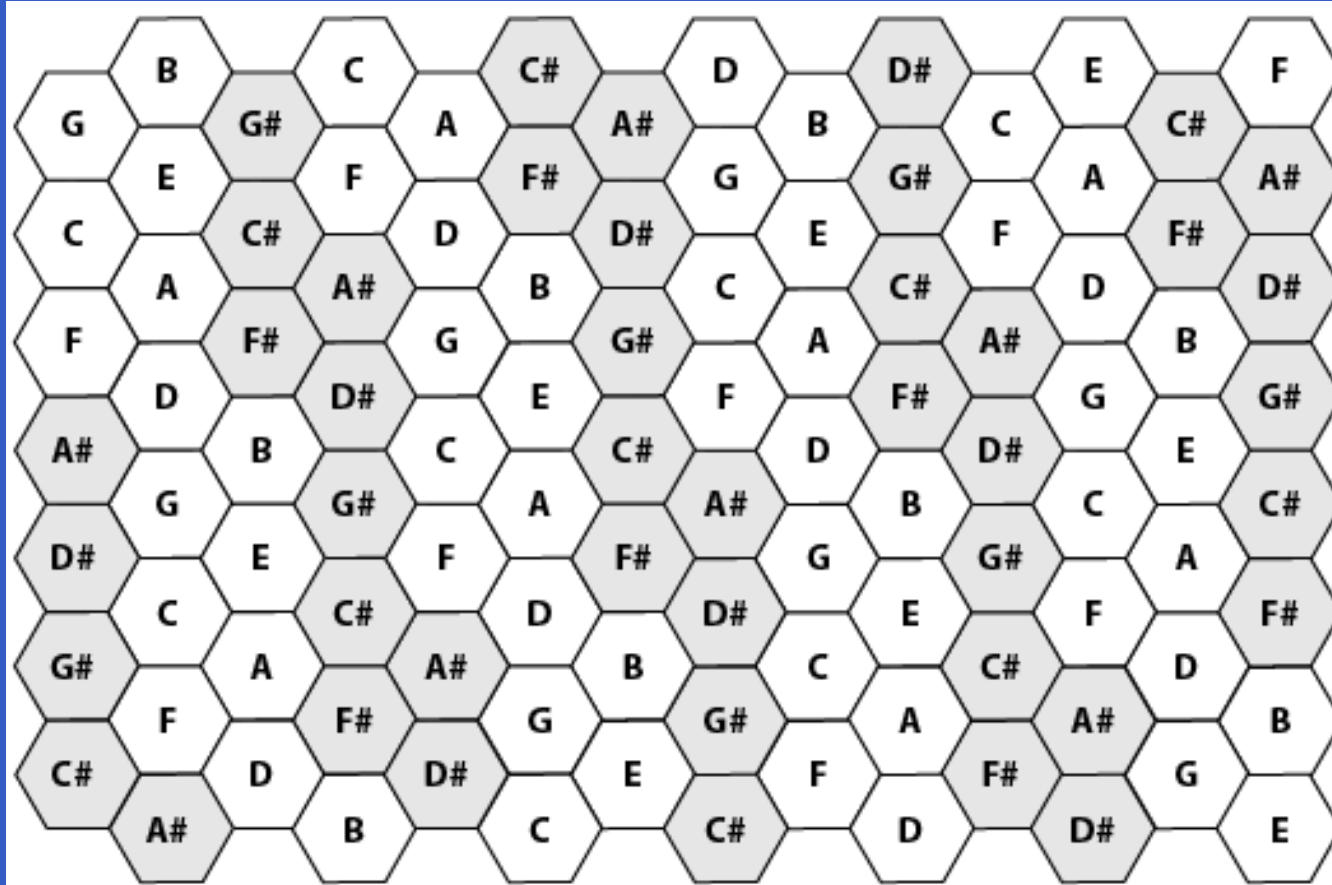
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The **structure** of the application should be such that it in principle is usable in a MIDI-studio setting.

The Harmonic Table



Running a Sample Configuration

The screenshot displays the Arpeggigon software interface. On the left, a large hexagonal grid is shown with several notes placed on it. The notes are represented by green hexagons and black circles with white symbols. The grid is labeled "Layer" in the top-left corner. On the right, a control panel is visible, featuring a tempo slider set to 120, volume and strength sliders, and various configuration options. The control panel includes a "Layer beat" dropdown set to "Quarter note", a "Beats per bar" dropdown set to 4, and a "Repeat count" section with checkboxes for "Enable repeat count" and "Keep heads on restart". The instrument is set to "Acoustic Grand Piano", and the note type is "Quarter note". The control panel also includes buttons for "Add layer", "Remove layer", "Save configuration", "Load configuration", "Restart", "Pause", "Stop", and "Record".

Layer

Tempo 120

Layer beat: Quarter note

Beats per bar: 4

Repeat count: 1

Acoustic Grand Piano

NoAccent

NoSlide

Quarter note

1

+ Add layer

- Remove layer

+ Save configuration

Load configuration

Restart

Pause

Stop

Record

The Rest of this Talk

- Brief introduction to FRP and Yampa
- The Arpeggigon core
- Brief introduction to Reactive Values and Relations
- The Arpeggigon shell

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Good conceptual fit for games, musical applications ...

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Yampa

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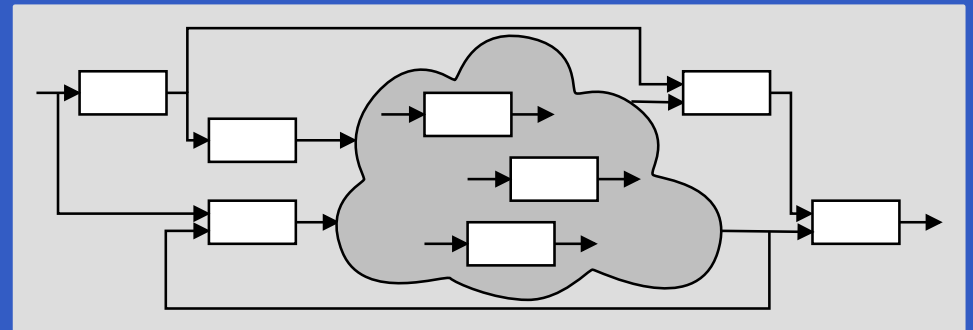
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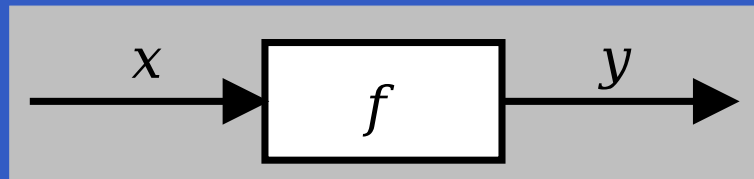
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 - **Signals**: time-varying values
 - **Signal Functions**: pure functions on signals
 - **Switching**: temporal composition of signal functions

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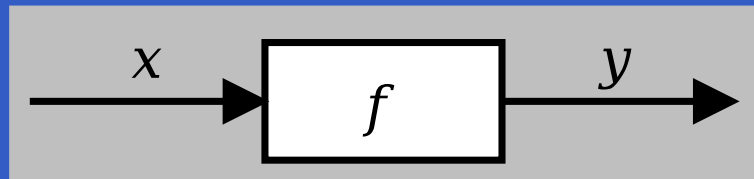
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- Programming model:



Signal Functions (1)

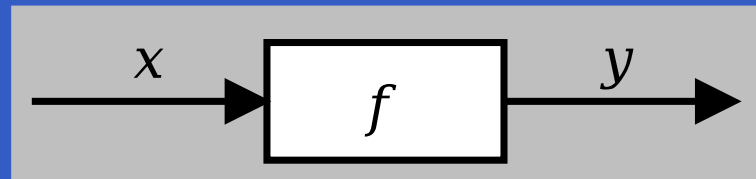


Signal Functions (1)



Intuition:

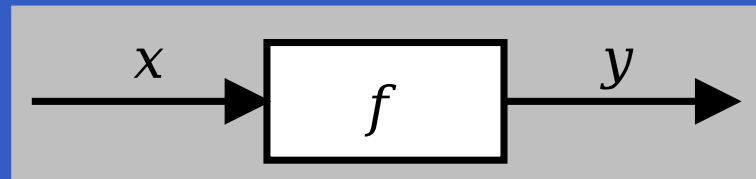
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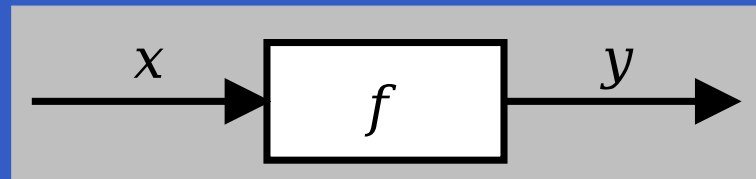
Time $\approx \mathbb{R}$

Signal $a \approx \text{Time} \rightarrow a$

$x :: \text{Signal } T1$

$y :: \text{Signal } T2$

Signal Functions (1)



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$Time \approx \mathbb{R}$

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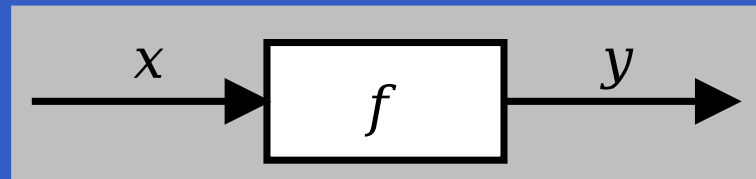
$x :: Signal\ T1$

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$SF\ a\ b \approx Signal\ a \rightarrow Signal\ b$

$f :: SF\ T1\ T2$

Signal Functions (1)



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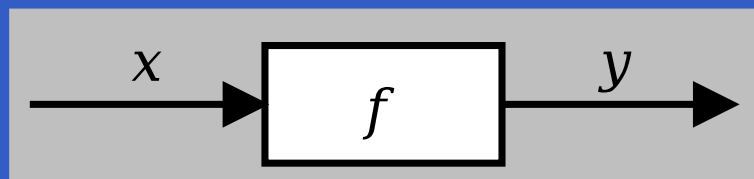
$y :: Signal\ T2$

$SF\ a\ b \approx Signal\ a \rightarrow Signal\ b$

$f :: SF\ T1\ T2$

Additionally, **causality** required: output at time t must be determined by input on interval $[0, t]$.

Signal Functions (2)



Example:

integral :: *VectorSpace* *a* *s* \Rightarrow *SF* *a* *a*

$$y(t) = \int_0^t x(\tau) d\tau$$

Clearly causal: output at time t determined by input on interval $[0, t]$.

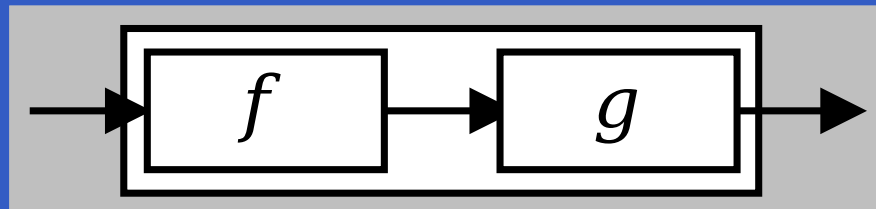
Composition

In Yampa, systems are described by combining signal functions (forming new signal functions).

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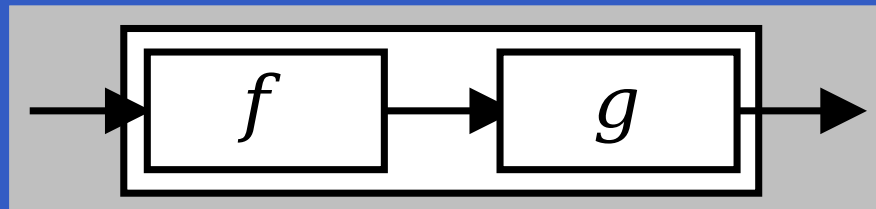
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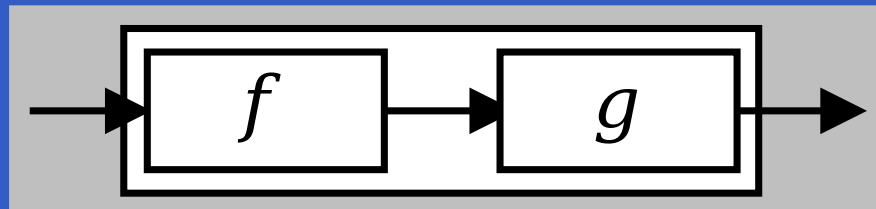
A **combinator** that captures this idea:

$$(\gg) :: SF\ a\ b \rightarrow SF\ b\ c \rightarrow SF\ a\ c$$

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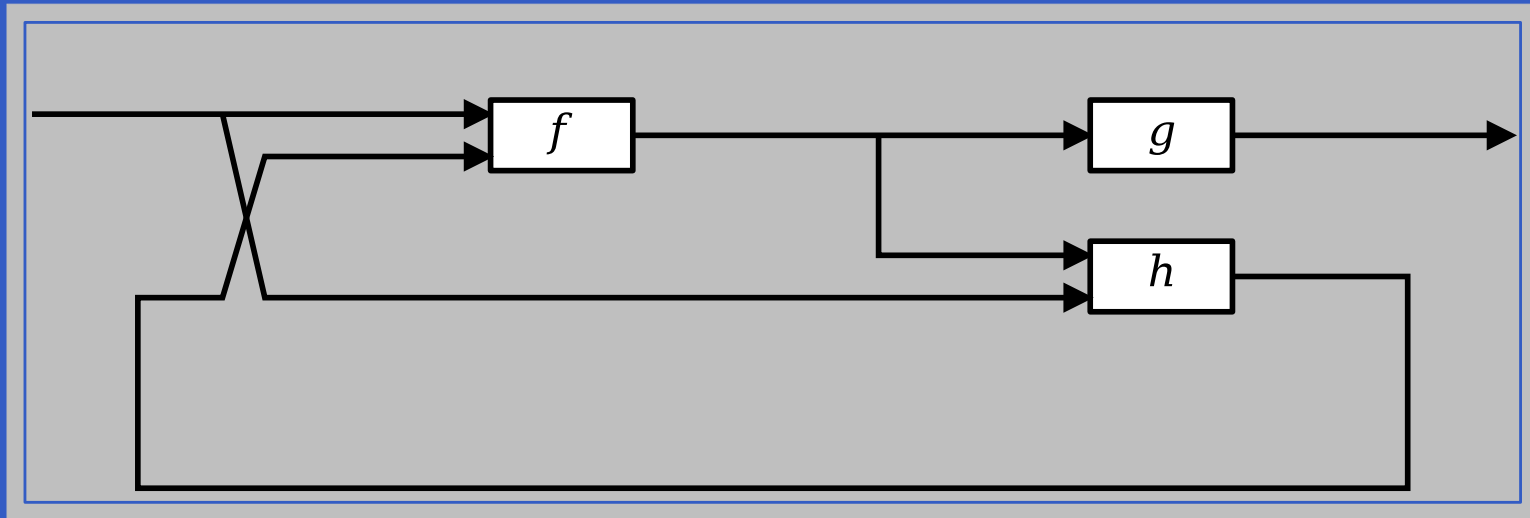


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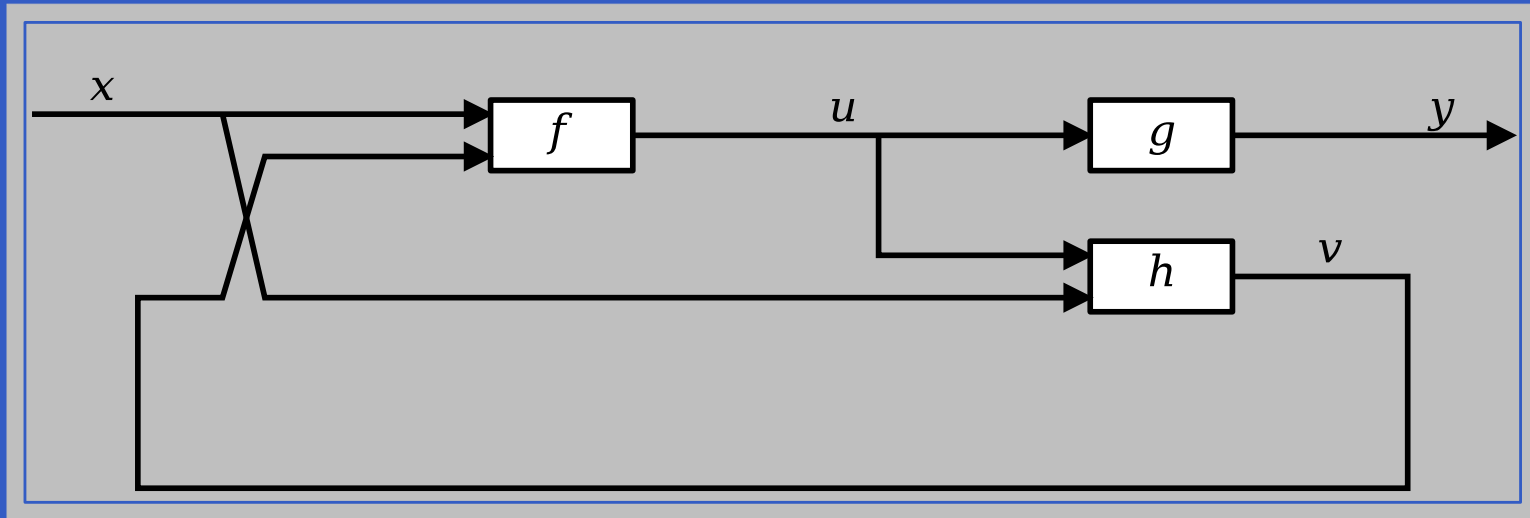
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Signal functions are the primary notion; signals a secondary one, only existing indirectly.

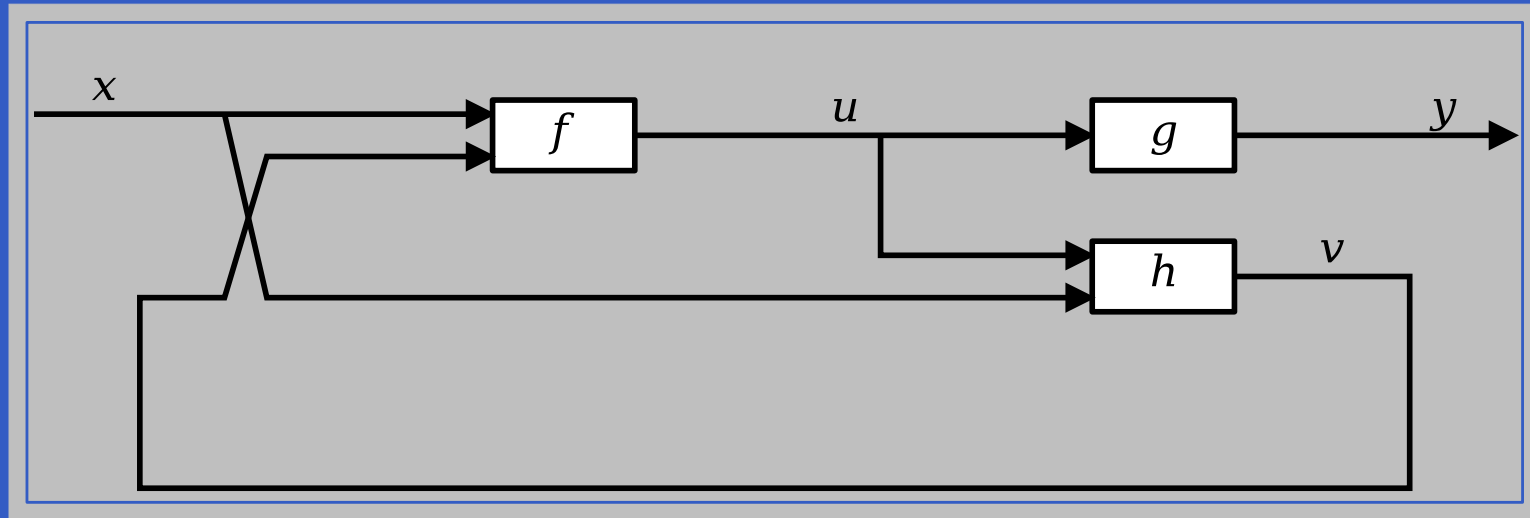
Arrow Notation



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`proc x → do`

`rec`

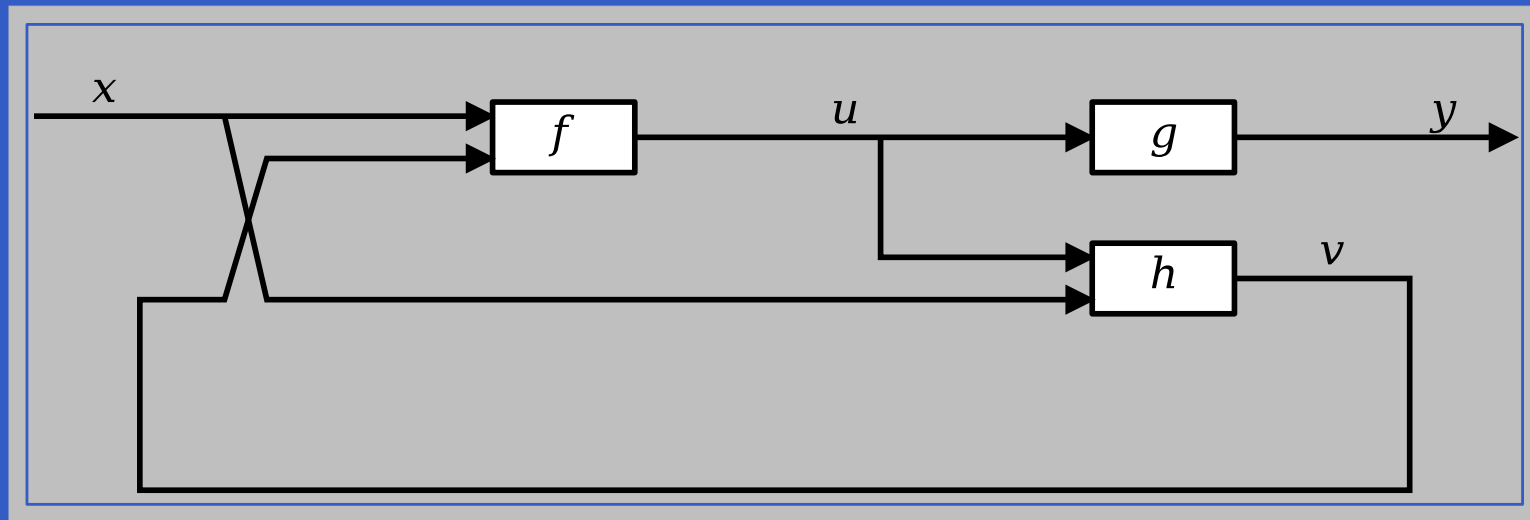
$u \leftarrow f \multimap (x, v)$

$y \leftarrow g \multimap u$

$v \leftarrow h \multimap (u, x)$

`return A \multimap y`

Arrow Notation



`proc x → do`

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Only syntactic sugar:
everything translated into a
combinator expression.

Events

Yampa models discrete-time signals by lifting the *co-domain* of signals using an option-type:

```
data Event a = NoEvent | Event a
```

Discrete-time signal = Signal (Event α).

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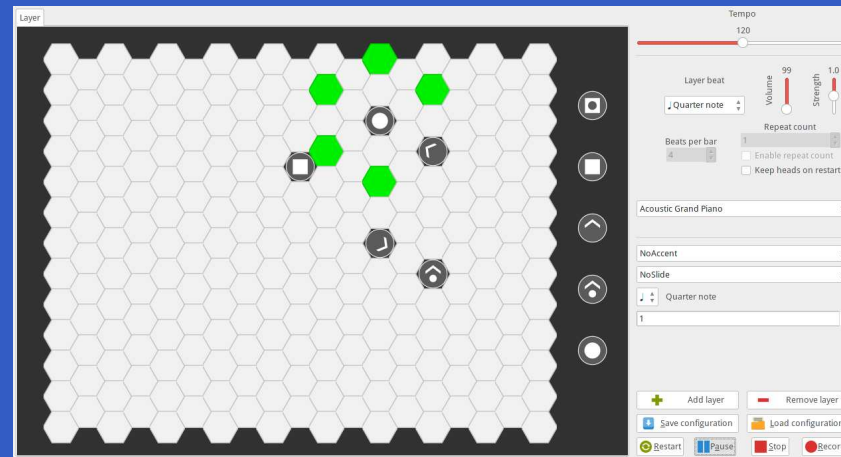
$$\text{data } \mathit{Event} \ a = \mathit{NoEvent} \mid \mathit{Event} \ a$$

Discrete-time signal = `Signal (Event a)`.

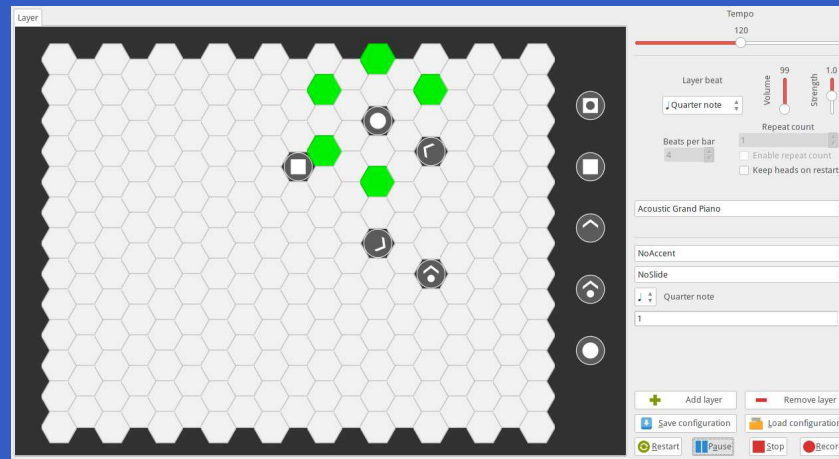
Some functions and event sources:

$$\mathit{tag} :: \mathit{Event} \ a \rightarrow b \rightarrow \mathit{Event} \ b$$
$$\mathit{after} :: \mathit{Time} \rightarrow b \rightarrow \mathit{SF} \ a \ (\mathit{Event} \ b)$$
$$\mathit{edge} :: \mathit{SF} \ \mathit{Bool} \ (\mathit{Event} \ ())$$

Aspects of the Arpeggigon

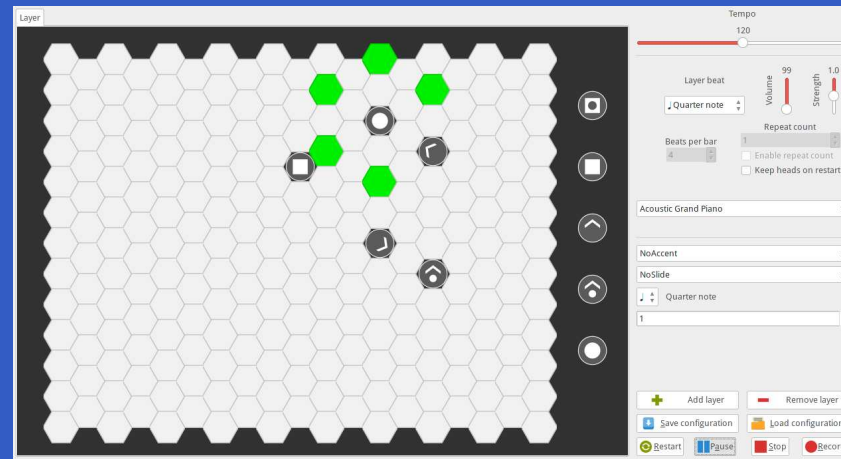


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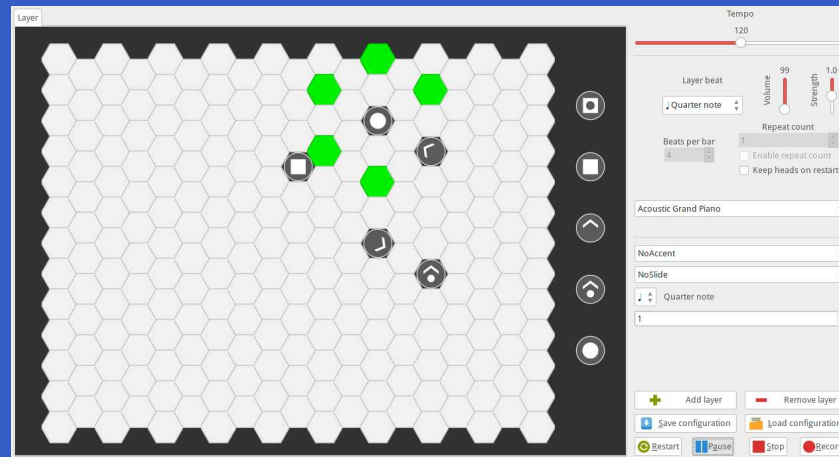
- *Interactive*

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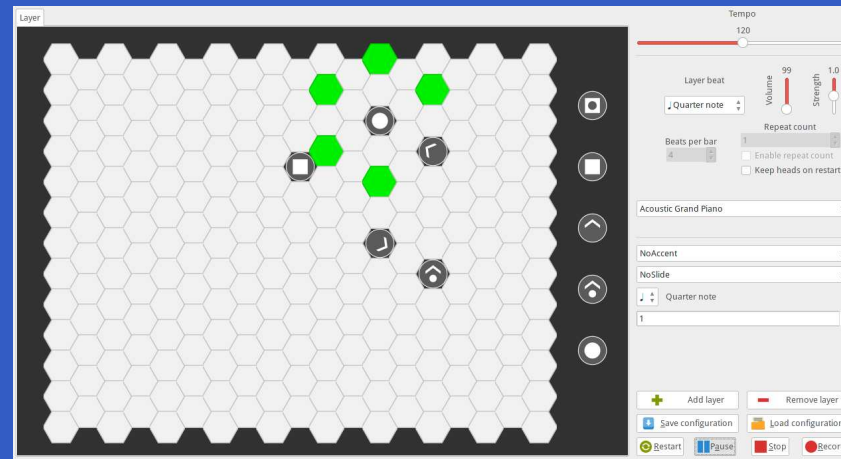
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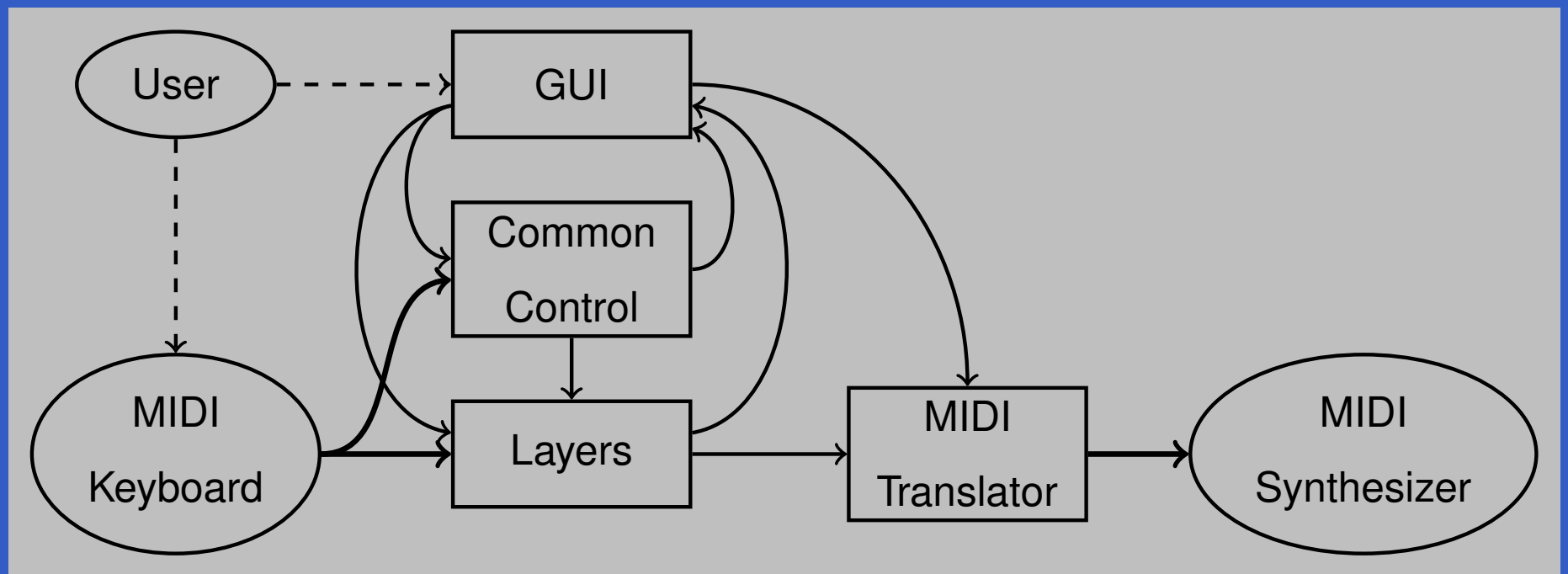
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Aspects of the Arpeggigon



- **Interactive**
- Layers can be added/removed: **dynamic structure**
- Notes generated at **discrete** points in time
- Configuration and performance parameters can be changed at **any** time

Arpeggigon Architecture



Cellular Automaton

State transition function for the cellular automaton:

$$\begin{aligned} advanceHeads &:: Board \rightarrow BeatNo \rightarrow RelPitch \rightarrow Strength \\ &\rightarrow [PlayHead] \rightarrow ([PlayHead], [Note]) \end{aligned}$$

Lifted into a signal function primarily using *accumBy*:

$$accumBy :: (b \rightarrow a \rightarrow b) \rightarrow b \rightarrow SF (Event a) (Event b)$$
$$\begin{aligned} automaton &:: [PlayHead] \\ &\rightarrow SF (Board, DynamicLayerCtrl, Event BeatNo) \\ &\quad (Event [Note], [PlayHead]) \end{aligned}$$

Automated Smooth Tempo Change

Smooth transition between two preset tempos:

$smoothTempo :: Tempo \rightarrow SF (Bool, Tempo, Tempo, Rate) \rightarrow Tempo$

$smoothTempo\ tpo0 = \mathbf{proc}\ (sel1, tpo1, tpo2, rate) \rightarrow \mathbf{do}$

\mathbf{rec}

$\mathbf{let}\ desTpo = \mathbf{if}\ sel1\ \mathbf{then}\ tpo1\ \mathbf{else}\ tpo2$

$diff = desTpo - curTpo$

$rate' = \mathbf{if}\ diff > 0.1\ \mathbf{then}\ rate$

$\mathbf{else\ if}\ diff < -0.1\ \mathbf{then}\ -rate$

$\mathbf{else}\ 0$

$curTpo \leftarrow arr\ (+tpo0) \lll integral \leftarrow rate'$

$\mathbf{return}\ A \leftarrow curTpo$

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 - GUI: GTK+
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- The Arpeggigon interacts with the outside world using two imperative toolkits:
 - GUI: GTK+
 - MIDI I/O: Jack
- Very imperative APIs: Hard or impossible to provide FRP wrappers.
- Instead, we use Ivan Perez's **Reactive Values and Relations** (RVR) to wrap the FRP core in a “shell” that acts as a bridge between the outside world and the pure FRP core.

Reactive Values and Relations (2)

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- RVs provide a uniform interface to GUI widgets, files, network devices, ...
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- RVR programming takes place in the IO monad, allowing arbitrary interfacing with imperative APIs.
- Yet, the high-level view is quite declarative/FRP-like.

System Tempo Slider

```
globalSettings :: IO (VBox, ReactiveFieldReadWrite IO Int)
globalSettings = do
  globalSettingsBox ← vBoxNew False 10
  tempoAdj          ← adjustmentNew 120 40 200 1 1 1
  tempoLabel       ← labelNew (Just "Tempo")
  boxPackStart globalSettingsBox tempoLabel PackNatural 0
  tempoScale      ← hScaleNew tempoAdj
  boxPackStart globalSettingsBox tempoScale PackNatural 0
  scaleSetDigits tempoScale 0
  let tempoRV =
        bijection (floor, fromIntegral)
        'liftRW' scaleValueReactive tempoScale
  return (globalSettingsBox, tempoRV)
```


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tempoRV' =

liftR2 (λtempo paused → if paused then 0 else tempo)

tempoRV

pauseButtonRV

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$$\begin{aligned} \text{tempoRV}' = & \\ & \text{liftR2 } (\lambda \text{tempo paused} \rightarrow \mathbf{if } \text{paused} \mathbf{ then } 0 \mathbf{ else } \text{tempo}) \\ & \text{tempoRV} \\ & \text{pauseButtonRV} \end{aligned}$$

- This is an equation defining $\text{tempoRV}'$ once and for all.

Connecting the Core to the Shell

The following function makes a signal function available as RVs:

yampaReactiveDual ::

a

→ SF a b

→ IO (ReactiveFieldWrite IO a, ReactiveFieldRead IO b)

This creates two reactive values: one for the input and one for the output of the signal function.

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- Reactive Values and Relations proved very helpful for bridging the gap between the outside world and the FRP core in a fairly declarative way.
- Performance in terms of overall execution time and space perfectly fine.
- **Timing** is not yet as tight as it should be due to naive MIDI generation.

Reading (1)

- Henrik Nilsson and Gueric Chupin. Funky Grooves: Declarative Programming of Full-Fledged Musical Applications. In *9th International Symposium on Practical Aspects of Declarative Languages (PADL 2017)*, pp. 163–172, January 2017.
- Ivan Perez and Henrik Nilsson. Bridging the GUI Gap with Reactive Values and Relations. In *Proceedings of the 8th ACM SIGPLAN Symposium on Haskell (Haskell'15)*, pp. 47–58, September 2015.

Reading (2)

- Henrik Nilsson, Antony Courtney, and John Peterson. Functional reactive programming, continued. In *Proceedings of the 2002 Haskell Workshop*, pp. 51–64, October 2002.
- Antony Courtney and Henrik Nilsson and John Peterson. The Yampa Arcade. In *Proceedings of the 2003 Haskell Workshop*, pp. 7–18, August 2003.