

CATEGORY THEORY  
MIDLANDS GRADUATE SCHOOL 2023

EXERCISE 1 (2 APRIL)

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REMINDER

**Definition 1** (category). A category  $\mathcal{C}$  consists of:

- a collection  $\mathcal{C}_0$  of objects, written  $X, Y, Z, \dots$
- for any two objects  $X$  and  $Y$ , a collection  $\mathcal{C}(X, Y)$  of morphisms, written  $f, g, h, \dots$
- a composition operation: for  $f \in \mathcal{C}(X, Y)$  and  $g \in \mathcal{C}(Y, Z)$ , we have  $g \circ f \in \mathcal{C}(X, Z)$
- for any object  $X$ , the identity morphism  $\text{id}_X \in \mathcal{C}(X, X)$

such that:

- Every identity morphism is left- and right-neutral. This means that, for  $f \in \mathcal{C}(X, Y)$ , we have  $f \circ \text{id}_X = f$  and  $\text{id}_Y \circ f = f$ .
- Composition is associative, i.e.  $(h \circ g) \circ f = h \circ (g \circ f)$ .

**Definition 2** (initial and terminal object). An object  $X$  in a category  $\mathcal{C}$  is initial if, for every object  $Y$ , there is exactly one morphism from  $X$  to  $Y$ . This means that  $\mathcal{C}(X, Y)$  is the one-element set. An object  $Z$  is terminal if, for every object  $Y$ , there is exactly one morphism from  $Y$  to  $Z$ , i.e.,  $\mathcal{C}(Y, Z)$  is the one-element set.

**Definition 3** (isomorphism). A morphism  $f \in \mathcal{C}(X, Y)$  is an isomorphism if there is a morphism  $g \in \mathcal{C}(Y, X)$  such that  $g \circ f = \text{id}_X$  and  $f \circ g = \text{id}_Y$ .

EXERCISE 1: THE FREE CATEGORY ON A DIRECTED MULTIGRAPH

A directed multigraph (sometimes also called *quiver*) consists of:

- a set  $V$  of vertices
- for each pair  $(a, b)$  of vertices, a set  $E(a, b)$  of edges from  $a$  to  $b$ .

Note that the edges are directed, i.e.,  $E(a, b)$  is not necessarily the same as  $E(b, a)$ , and many parallel edges are allowed.

Let  $G = (V, E)$  be a directed multigraph. The *free category on  $G$* , (here) written  $\mathcal{F}_G$ , has  $V$  as objects and a morphism  $\mathcal{F}_G(X, Y)$  is a sequence of consecutive edges, starting in  $X$  and ending in  $Y$ .

- a. Show that  $\mathcal{F}_G$  is a category.
- b. What are the isomorphisms in  $\mathcal{F}_G$ ?
- c. For which  $G$  does  $\mathcal{F}_G$  have an initial object? And when does  $\mathcal{F}_G$  have a terminal object?

EXERCISE 2: THE CATEGORY REL

The objects of the category REL are sets. For sets  $X$  and  $Y$ , an element of  $\text{REL}(X, Y)$  is a *relation* between  $X$  and  $Y$ , i.e., a subset  $R \subseteq (X \times Y)$ . The composition of relations  $R \subseteq (X \times Y)$  and  $S \subseteq (Y \times Z)$  is given by

$$\{ (x, z) \mid \exists y \in Y. (x, y) \in R \wedge (y, z) \in S \}.$$

- a. Can you define identities and prove that REL is a category?
- b. What are the initial and terminal object of REL?
- c. What are the isomorphisms?

## BONUS EXERCISE: THE (ALMOST-)CATEGORY SPAN

The objects of **SPAN** are sets. A morphism between sets  $X$  and  $Y$  consists of a set  $Z$  together with a function  $f : Z \rightarrow X$  and a function  $g : Z \rightarrow Y$ .

- a.** The composition operation works in a similar way as for **REL**, but we don't need an existential quantifier this time. Can you make this operation precise?
- b.** Can you prove associativity? Can you define identities and prove that **SPAN** is a category? What works or fails? (Hint: It only almost works. **SPAN** is a structure that is slightly weaker than a category, namely a bicategory.)
- c.** What are the initial and terminal object of **SPAN**?
- d.** What are the isomorphisms?
- e.** How does **SPAN** compare to the category **REL**?