

Decidability and Semidecidability via Ordinals

Nicolai Kraus

jww **Fredrik Nordvall Forsberg** and **Chuangjie Xu**

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What are ordinal numbers?

One answer: **Numbers** for counting/ordering:

0, 1, 2, 3, ... ω , $\omega + 1$, $\omega + 2$, ... $\omega \cdot 2 + 19$, ...

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Builds & sells houses



Bob

Wants to buy

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20 days, at most!

When is my house ready?



Alice

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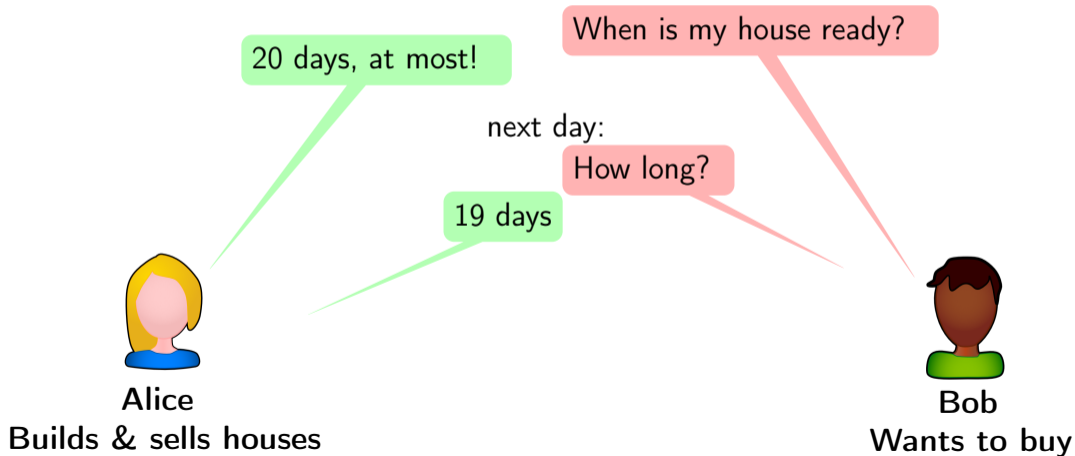
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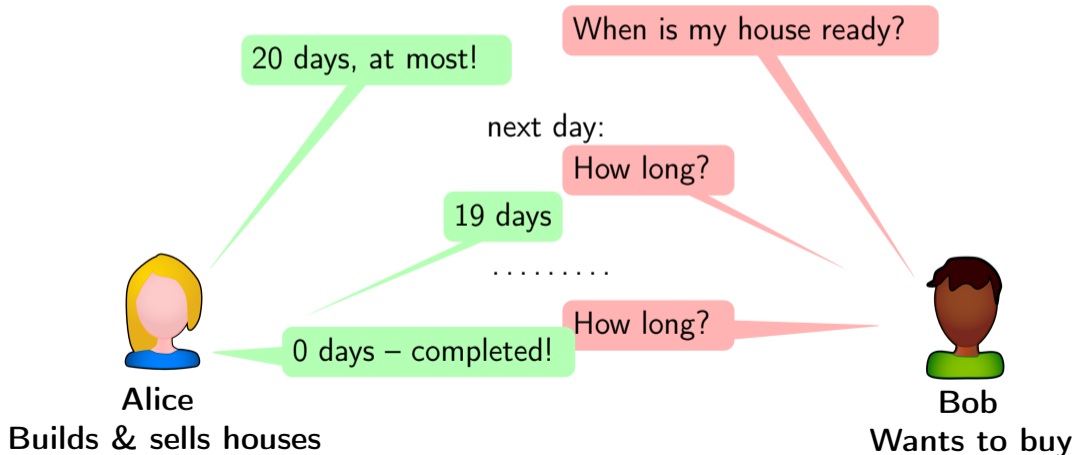
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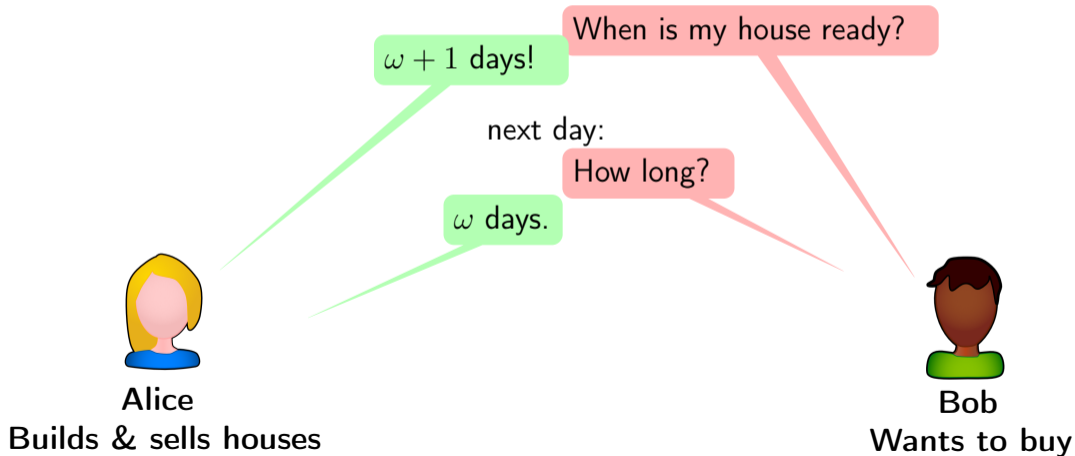
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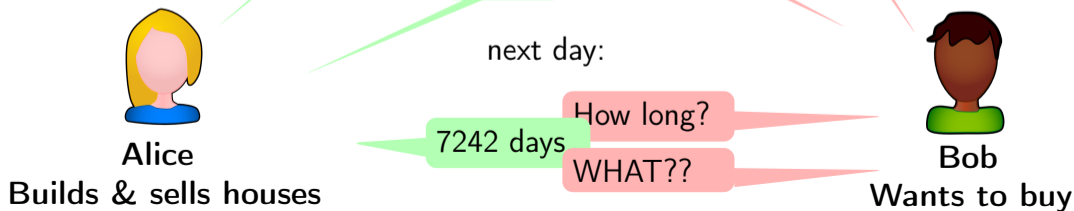
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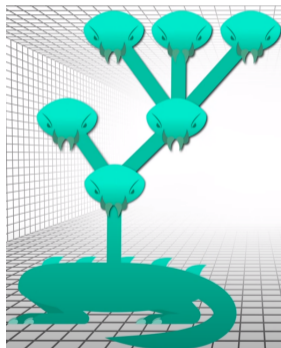
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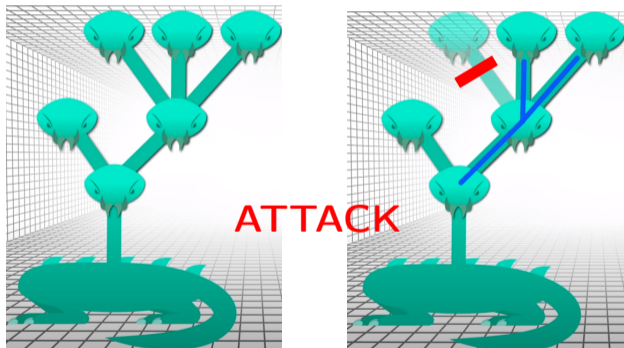


Hydra by Kirby and Paris 1982, and pictures by PBS Infinite Series, <https://youtu.be/uWwUpEY4c8o>

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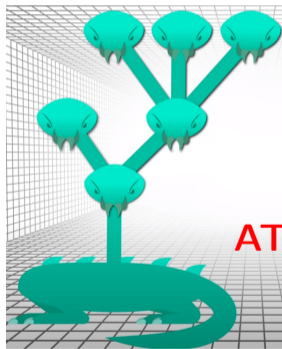


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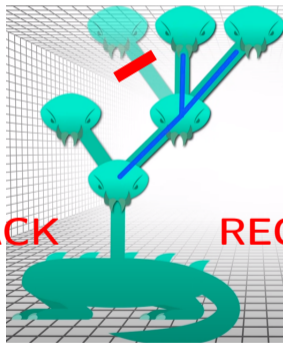
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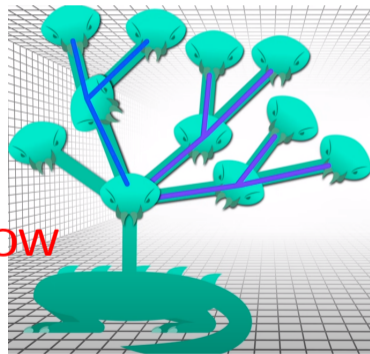
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ATTACK



REGROW



Brouwer ordinal trees in constructive type theory

Inductive type \mathcal{B} of Brouwer trees: data \mathcal{B} where

- zero : \mathcal{B}
- succ : $\mathcal{B} \rightarrow \mathcal{B}$
- limit : $(\mathbb{N} \rightarrow \mathcal{B}) \rightarrow \mathcal{B}$

Then: Define $\omega := \text{limit}(0, 1, 2, 3, \dots)$
 $\omega \cdot 2 := \text{limit}(\omega, \omega + 1, \omega + 2, \dots)$
and so on (addition, multiplication, exponentiation are standard).

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One problem (for our application): **limit**(0, 1, 2, 3, ...) \neq **limit**(1, 2, 3, ...)

Our approach: induction-induction and path constructors, ensuring:

- ▶ Limits can only be taken of strictly increasing sequences;
- ▶ Bisimilar sequences have equal limits.

In cubical Agda:

```
data Brw where
  zero  : Brw
  succ  : Brw → Brw
  limit : (f : ℕ → Brw) → {f↑ : increasing f} → Brw
  bisim : ∀ f {f↑} g {g↑} →
    f ≈ g →
    limit f {f↑} ≡ limit g {g↑}
  trunc : isSet Brw
```

note: $x < y$
means $\text{succ } x \leq y$

```
data _≤_ where
  ≤-zero      : ∀ {x} → zero ≤ x
  ≤-trans     : ∀ {x y z} → x ≤ y → y ≤ z → x ≤ z
  ≤-succ-mono : ∀ {x y} → x ≤ y → succ x ≤ succ y
  ≤-cocone    : ∀ {x} f {f↑ k} → (x ≤ f k) → (x ≤ limit f {f↑})
  ≤-limiting  : ∀ f {f↑ x} → ((k : ℕ) → f k ≤ x) → limit f {f↑} ≤ x
  ≤-trunc     : ∀ {x y} → isProp (x ≤ y)
```

Everything that one can “reasonably expect” works: $<$ is wellfounded, \leq is anti-symmetric, limits are actually limits, arithmetic operations work, and so on.

Decidability properties

P is *decidable* if we can prove $P \uplus \neg P$.

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Sure: No for zero, yes for limits; for succ y , check whether $y > 41$.

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4. $x > \omega$?

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4. $x > \omega$?

Can decide it for zero and succ, but: $\text{limit}(x_0, x_1, x_2, \dots) > \omega$?

When is $\text{limit}(x_0, x_1, x_2, \dots) > \omega$?

- ▶ For any i , we can check whether x_i is finite.
- ▶ As soon as we discover an infinite x_i , the question is decided positively.
- ▶ Only if all x_i are finite, the answer is negative.

\Rightarrow *Semidecidable*.

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\Rightarrow *Semidecidable*.

Definition (Bauer 2006): P is semidecidable if

$$\exists (s : \mathbb{N} \rightarrow \text{Bool}). P \leftrightarrow \exists k. s_k = \text{true}$$

(Note: $\exists (x : A). B(x)$ means $\|\Sigma(x : A). B(x)\|$.)

Fact: For any x , the question $x > \omega$ is semidecidable.

The other direction

Given $s : \mathbb{N} \rightarrow \text{Bool}$, we can construct an increasing sequence f by:

$$f\ 0 \equiv \text{zero}$$

$$f\ (n + 1) = \begin{cases} (f\ n) + \omega & \text{if } n \text{ is [minimal] such that } s_n = \text{true} \\ \text{succ}(f\ n) & \text{else.} \end{cases}$$

Then: $(\text{limit } f > \omega) \leftrightarrow (\exists k. s_k = \text{true})$.

Semidecidability via ordinals

Via these translations: For any proposition P ,

$$\exists(y : \mathcal{B}). P \leftrightarrow (y > \omega) \quad \longleftrightarrow \quad \exists(s : \mathbb{N} \rightarrow \text{Bool}). P \leftrightarrow \exists k. s_k = \text{true}$$

“ P decidable in ω steps” (??)

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“ P decidable in ω steps” (??)

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What if we swap ω for another ordinal α ?

$$\exists(y : \mathcal{B}).P \leftrightarrow (y > \alpha) \quad \text{“decidable in } \alpha \text{ steps”}$$

(or $y \geq \alpha$, $y = \alpha$, any $Q(y)$, ...)

Fewer than ω steps

Let n be a natural number. Then:

$$\begin{array}{ccc} \exists(y : \mathcal{B}).P \leftrightarrow (y > n) & \longleftrightarrow & P \uplus \neg P \\ \text{"}P \text{ decidable in } n \text{ steps"} & & \text{"}P \text{ decidable"} \end{array}$$

More than ω steps – an example

Twin prime conjecture (TPC):

There are arbitrarily large numbers p such that p and $p + 2$ are both prime.

It's clearly semidecidable whether there is a twin pair $> 10^{1,000,000}$, but TPC doesn't seem to be semidecidable.

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It's clearly semidecidable whether there is a twin pair $> 10^{1,000,000}$, but TPC doesn't seem to be semidecidable.

However, one can show:

$$\exists(y : \mathcal{B}).\text{TPC} \leftrightarrow (y = \omega^2)$$

“TPC is decidable in ω^2 steps.”

(Note: $=$ can be replaced by $>$ or \geq .)

TPC's ordinal

Define a sequence $f : \mathbb{N} \rightarrow \mathcal{B}$ by:

$$f 0 \equiv \text{zero}$$

$$f(n+1) = \begin{cases} (f n) + \omega & \text{if } n \text{ and } n+2 \text{ are prime} \\ (f n) + 1 & \text{else.} \end{cases}$$

Claim: $\text{TPC} \leftrightarrow \text{limit } f = \omega^2$ [$\leftrightarrow \text{succ}(\text{limit } f) > \omega^2$]

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Sketch $\text{TPC} \rightarrow (\text{limit } f = \omega^2)$:

For any n , we find k s.t. there are at least n twin prime pairs below k , thus

$f_k \geq \omega \cdot k$, thus $\text{limit } f \geq \omega \cdot \omega$.

At the same time, f never exceeds ω^2 .

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Sketch $(\text{limit } f \geq \omega^2) \rightarrow \text{TPC}$:

- $(\text{limit } f \geq \omega^2) \Rightarrow \exists k. f_k \geq \omega \cdot n$
- $\Rightarrow \exists k. \neg\neg(\text{There are at least } n \text{ twin primes } \leq k)$
- $\Rightarrow \exists k. \text{There are at least } n \text{ twin primes } \leq k$
- $\Rightarrow \Sigma k : \mathbb{N}. \text{There are at least } n \text{ twin primes } \leq k$
- $\Rightarrow \text{There is a twin prime pair above } n$

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Claim: TPC \leftrightarrow limit $f = \omega^2$ [\leftrightarrow succ(limit f) $> \omega^2$]

Thanks for your attention!

- $(\text{limit } f \geq \omega^2) \Rightarrow \exists k. f_k \geq \omega \cdot n$
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