

Computer Systems Architecture

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Lecture 11: Pointers and References



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What does the following C program print?

```
void swap(int x,int y) {
    int z;
    z = x;
    x = y;
    y = z;
}

int main() {
    int a,b;
    a = 2;
    b = 3;
    swap(a,b);
    printf("a=%d, b=%d\n",a,b);
}
```



References in C

- We can declare pointer types in C, e.g.

```
int *x;
```

means that `x` holds a pointer to an integer.

- To dereference a pointer we also use `*`, e.g. `*x` has type `int`.

- The operator `&` returns a pointer to a variable.

- E.g. if we have declared

```
int y
```

then `&y` has type `int *`, *pointer to an integer*.



What does the following C program print?

```
void swap(int *x,int *y) {
    int z;
    z = *x;
    *x = *y;
    *y = z;
}

int main() {
    int a,b;
    a = 2;
    b = 3;
    swap(&a,&b);
    printf("a=%d, b=%d\n",a,b);
}
```



What about Java?

- Java hasn't got pointer types.
- Basic datatypes are always passed *by value*.
- Objects, arrays and strings are passed as references.
- Java avoids pointer bugs, which are common and hard to detect.



What does the following Java program print?

```
class Int {
    int n;
    Int(int m) { n = m; } }

public class Swap {
    static void swap(Int x, Int y) {
        int z;
        z = x.n;
        x.n = y.n;
        y.n = z; }
    public static void main(String args[] ) {
        Int a = new Int(1);
        Int b = new Int(2);
        swap(a,b);
        System.out.println("a="+a.n+" b="+b.n); }
}
```



swap in MIPS

```

        .data
aa:     .word 1
bb:     .word 2

        .text
        .globl main
main:   la $a0, aa
        la $a1, bb
        jal swap          # swap(&a,&b);
        ...              # print a,b

swap:   # x=$a0, y=$a1, z=$t0
        lw $t0,($a0)     # z = *x;
        lw $t1,($a1)
        sw $t1,($a0)     # *x = *y;
        sw $t0,($a1)     # *y = z;
        jr $ra

```



Arrays

- One of the most basic data structures in CS
- Usually a block of consecutive elements in memory
 - All same size (s bytes); same offset from one to the next
 - The i^{th} element is at offset $i \times s$ bytes from beginning
 - Looking up an element of the array is termed '*indexing*'
- Characterised by constant-time indexing
 - No more faster to look up `xs[0]` than `xs[42]`
 - Contrast this with a linked-list¹ (not in this course)
- We can implement arrays using pointer arithmetic
- e.g. Assembly equivalent of an `int[]` in Java/C would be...
 - a consecutive block of word-sized signed integers
 - represented by its starting address and length

¹like lists in Haskell



Using (Integer) Arrays: C

```
int as[8] = { 3, 1, 4, 1, 5, 9, 2, 6 };
```

```
int array_max(int xs[], int length) {  
    int i,m;  
    m = INT_MIN;  
    for(i=0; i<length; i++) {  
        if(m < xs[i])  
            m = xs[i];  
    }  
    return m;  
}
```

```
int main() {  
    printf("max = %d\n", array_max(as,8));  
}
```



Using (Integer) Arrays: Assembly, Part 1

```

array_max: # $a0: array address, $a1: array length
    li $v0, 0x80000000    # MIN_VALUE
    li $t0, 0             # i = 0
    j am_cond

am_loop:
    sll $t1, $t0, 2       # 4 * i
    add $t1, $a0, $t1     # xs + 4*i bytes
    lw $t1, ($t1)         # lookup xs[i]
    addi $t0, $t0, 1      # i++
    bge $v0, $t1, am_cond # if(m < xs[i])
        move $v0, $t1    # m = xs[i]

am_cond:
    blt $t0, $a1, am_loop # i < length?
    jr $ra

```



Using (Integer) Arrays: Assembly, Part 2

```
.data
as:  .word 3, 1, 4, 1, 5, 9, 2, 6
.text
.globl main
main: addi $sp, $sp, -4
      sw $ra, 0($sp)
      la $a0, as      # $a0 = address of as
      li $a1, 8       # $a1 = as.length
      jal array_max  # array_max(as, as.length)
      move $a0, $v0
      li $v0, 1       # print_int
      syscall
      lw $ra, 0($sp)
      addi $sp, $sp, 4
      jr $ra
```



Strings

- Java strings are opaque objects of class `String`
- Assembly strings are arrays of ASCII characters, or bytes
 - End marked with a NUL, rather than storing its length
- You've already used them before
 - with the `.asciiz` directive
 - and the `print_string` syscall
- What else can we do with strings?



String length in C

```
int strlen(char *s) {
    int l;
    l = 0;
    while(*s != 0) {
        s++;
        l++;
    }
    return l;
}

int main() {
    printf("%d\n",strlen("hello"));
}
```



String Length in Assembler

```
strlen: # s=$a0,l=$v0
li $v0, 0          # l = 0 ;
j strlen_cond
strlen_loop:
addi $v0, $v0, 1   # l++
strlen_cond:
lbu $t0, ($a0)
addi $a0, $a0, 1   # s++
bne $t0, $zero, strlen_loop # while(*s != '\0')
jr $ra
```



String Length in Assembler

```
        .data
hello:  .ascii "hello"
        .text
        .globl main
main:   la $a0,hello
        jal strlen
        move $a0,$v0
        li $v0, 1      # print_int
        syscall
        li $v0,10
        syscall        # exit
```



strcat, 1st attempt

```
char* strcat(char *s, char *t) {
    char *r;
    r = s;
    while(*s != '\0') s++;
    do {
        *s = *t;
        s++;
        t++;
    } while(*t != '\0');
    return r;
}
```

```
int main() {
    printf("%s\n", strcat("hello ", "world"));
}
```



Oops!

```
sean:code txa$ strcat1  
Bus error
```



Dynamic Data

- So far we've only dealt with *static* data
 - Contents may change, but size and location doesn't
 - Same sense as the `static` keyword in Java
- In Java, "hello" + "world" concatenates two strings
 - But neither of the original strings are modified
 - Instead a new string is created on the *heap*
- The heap is a much larger pool of memory than the stack
 - In C we can allocate data using `malloc`
 - Unused data can be returned by using `free`
- Storage allocated on the heap persist across procedures
 - Caller can't access stack storage



strcat, 2nd attempt

```
char* strcat(char *s, char *t) {
    char *r,*u;
    r = (char *) malloc(strlen(s)+strlen(t)+1);
    u = r;
    while(*s != '\0') {
        *u = *s;
        s++;
        u++;
    }
    do {
        *u = *t;
        u++;
        t++;
    } while(*t != '\0');
    return r;
}
```



Horrors of Memory Leaks

```
int main() {
    char *s;
    while(1) {
        s = malloc(1000);
        *s='x';
        printf(".");
    }
}
```



Horrors of Memory Leaks

- Program uses up all memory and will eventually crash.
- Small leaks hard to discover: may run for a long time



Java version

```
public class Foo {  
    public static void main(String[] args) {  
        while(true) {  
            int[] as = new char[1000];  
            as[0] = 'x';  
        }  
    }  
}
```



Automatic Garbage Collection

- Java has automatic garbage collection
 - Inaccessible objects are periodically freed by JVM
 - SPIM doesn't/can't have automatic garbage collection
- Can you write a Java program which runs out of memory?

