G52MAL Machines and Their Languages Lecture 13

Pushdown Automata (PDA)

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 - Through automata: a mechanical procedure for deciding whether or not a word belongs to the language.
 - Through *regular expressions*: a direct description of the words in a language.
- We have learned that not all languages are regular. In particular, the class of context-free languages includes the regular languages but is strictly larger.

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In other words: is there a mechanical procedure for determining language membership for the context-free languages, and if so, what is the "simplest" such machine?

The answer is YES! Pushdown Automata (PDA).

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 - Suffices for describing the context-free languages (CFL).
 - Is not enough to describe languages more general than the CFLs (as we will see).