

Introduction to FRP and Yampa through Games and Music

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Aims and Overview (1)

- Introduction to *Functional* Reactive Programming

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- Demonstrating that a pure, declarative, reactive approach covers more applications areas than what one might think.

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- Introduction to **Functional** Reactive Programming
- Demonstrating that a pure, declarative, reactive approach covers more applications areas than what one might think.
- Through the above, give you a hopefully somewhat different and useful perspective on the reactive programming technologies that you are already using.

Aims and Overview (2)

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- Circuit-like programming metaphor.
- Transforming streams by e.g. mapping or accumulation.
- Similar primitive streams, e.g. never, periodic.

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Aims and Overview (3)

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- Time (notionally) **continuous**.
- But **events** (discrete time), also supported, allowing for **hybrid** systems.
- Synchronous (no race conditions)
- Declarative (emphasis on what, not how; fewer operational concerns)

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Aims and Overview (4)

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- High-level support for highly dynamic system structure (generalisation of XStream’s `flatten`, no need for low-level attaching/detaching of listeners).

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And specific to Yampa:

- “Stream processors” is the central, first-class, abstraction, while “streams” are secondary (very close to the circuit metaphor).
- High-level support for highly dynamic system structure (generalisation of XStream’s `flatten`, no need for low-level attaching/detaching of listeners).
- Statically typed

Example: Feedback

Arranging feedback in XStream necessitates taking operational concerns, *how*, into account:

```
var secondProxy$ = xs.create();  
var first$ = secondProxy$.map ...;  
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```

Of course, the intent, *what*, is really:

```
var first$ = second$.map ...;  
var second$ = first$.map ...;
```

with “=” denoting equality rather than assignment.



Why Program Games Declaratively?

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Perhaps not so surprising:

- Many pragmatical reasons: performance, legacy issues, . . .
- State and effects are pervasive in video games: Is declarative programming even a conceptually good fit?

But Why **NOT**, Really?

Many eloquent and compelling cases for functional programming in general:

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One key point: Program with whole values, not a word-at-a-time. Which, of course, is the point also of frameworks like Cycle.js and XStreams.

Possible Gains (1)

With his Keera Studios hat on, Ivan's top three reasons:

- Reliability.
- Lower long-term maintenance cost.
- Lower production cost and fast time-to-prototype.

Possible Gains (2)

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The Unreal Engine

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The Unreal Engine

E.g. pure, declarative code:

- promotes parallelism
- eliminates many sources of errors

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“Whole Values” for Games?

How should we go about writing video games “declaratively”?

In particular, what should those “whole values” be?

- Could be conventional entities like vectors, arrays, lists and aggregates of such.
- Could even be things like pictures.

But we are going to go one step further and consider programming with *time-varying entities*.

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Functional Reactive Programming

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- FRP has evolved in a number of directions and into different concrete implementations.
- We will use Yampa: an FRP system embedded in Haskell.

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Our whole values are things like:

- The totality of input from the player
- The animated graphics output
- The entire life of a game object

We construct and work with *pure* functions on these:

- The game: function from input to animation
- In the game: fixed point of function on collection of game objects

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Take-home Message # 2

You too can program games declaratively ...

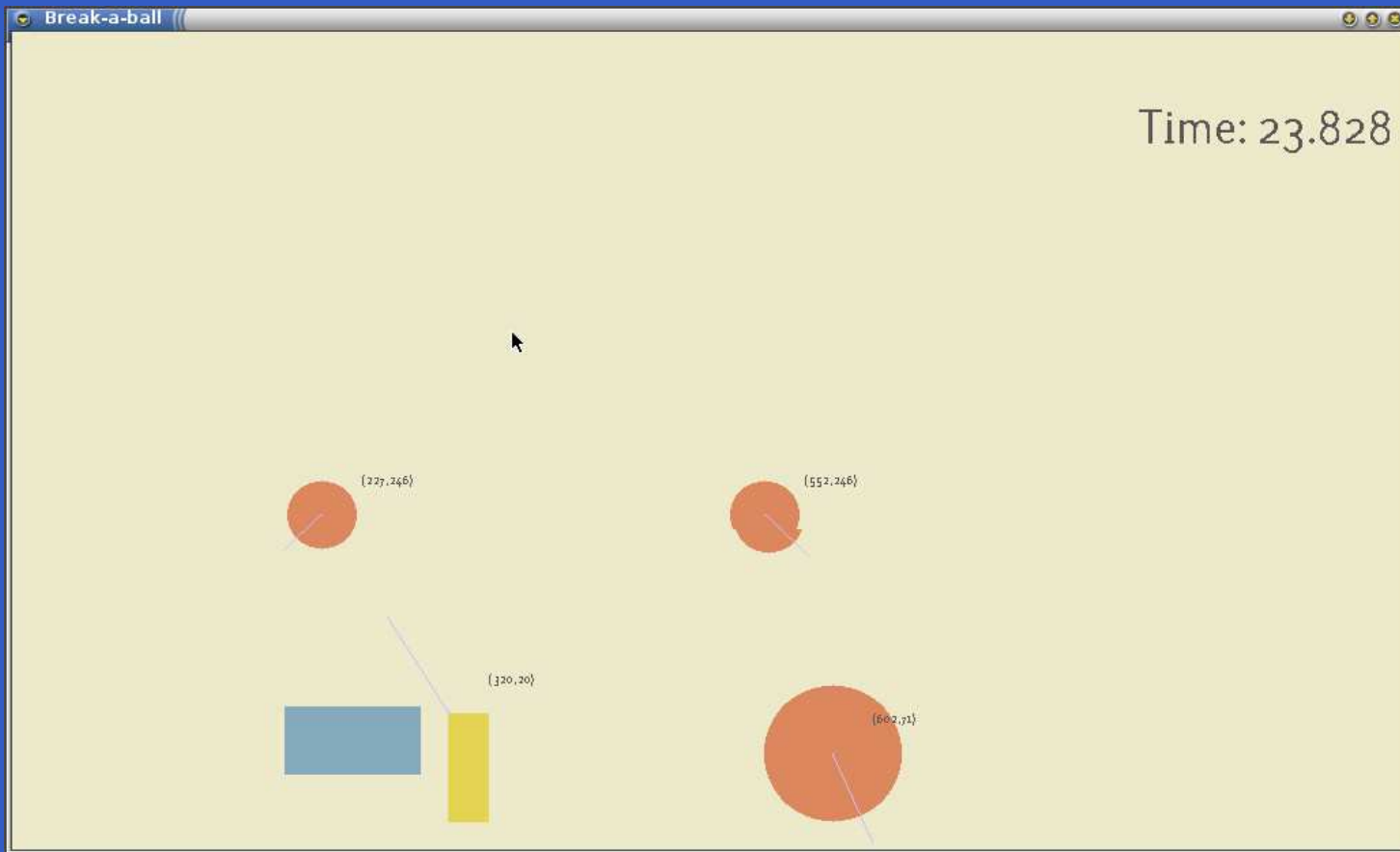
Take-home Message # 2

You too can program games declaratively ... today!



Take-home Game!

Or download one for free to your Android device!



Play Store: Pang-a-lambda (Keera Studios)

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Yampa

Yampa

- FRP implementation embedded in Haskell

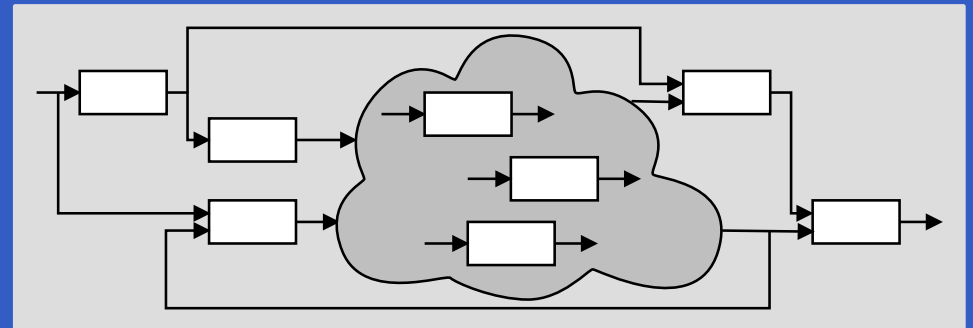
Yampa

- FRP implementation embedded in Haskell
- Key notions:
 - **Signals**: time-varying values (cf. streams)
 - **Signal Functions**: pure functions on signals
 - **Switching**: temporal composition of signal functions (cf. XStream's `flatten`)

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- Programming model:



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Yampa?

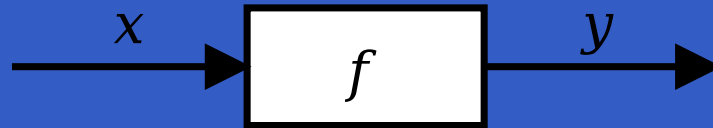
Yampa?

Yampa is a river with long calmly flowing sections and abrupt whitewater transitions in between.

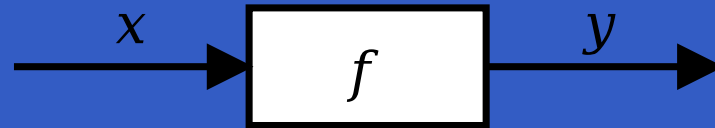


A good metaphor for hybrid systems!

Signal Functions

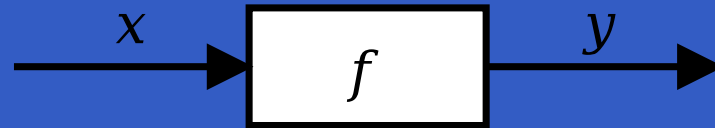


Signal Functions



Intuition:

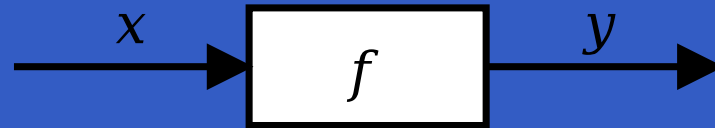
Signal Functions



Intuition:

$$Time \approx \mathbb{R}$$

Signal Functions



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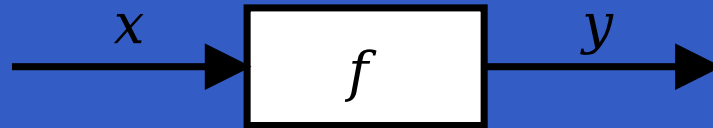
Time $\approx \mathbb{R}$

Signal $a \approx \text{Time} \rightarrow a$

$x :: \text{Signal } T1$

$y :: \text{Signal } T2$

Signal Functions



Intuition:

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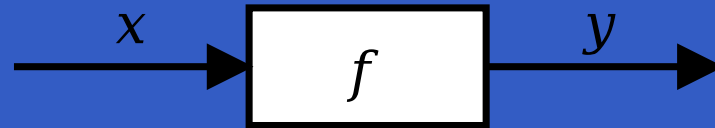
$x :: Signal\ T1$

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$SF\ a\ b \approx Signal\ a \rightarrow Signal\ b$

$f :: SF\ T1\ T2$

Signal Functions



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$x :: Signal\ T1$

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Additionally, **causality** required: output at time t must be determined by input on interval $[0, t]$.

Some Basic Signal Functions

identity :: SF a a

Some Basic Signal Functions

$identity :: SF\ a\ a$

$constant :: b \rightarrow SF\ a\ b$

Some Basic Signal Functions

identity :: $SF\ a\ a$

constant :: $b \rightarrow SF\ a\ b$

integral :: $VectorSpace\ a\ s \Rightarrow SF\ a\ a$

$$y(t) = \int_0^t x(\tau) d\tau$$

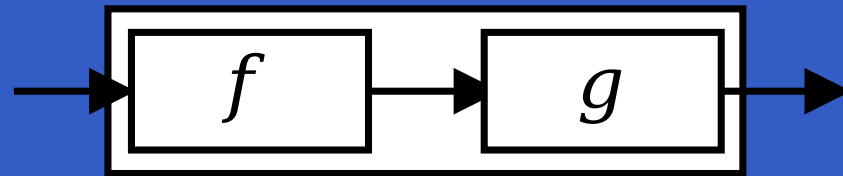
Composition

In Yampa, systems are described by combining signal functions (forming new signal functions).

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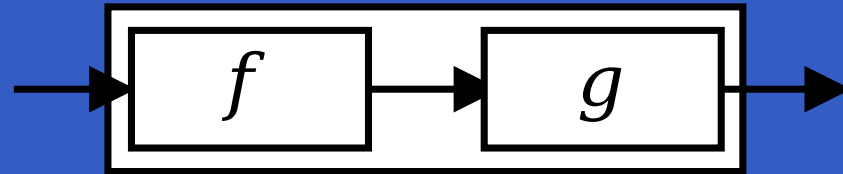
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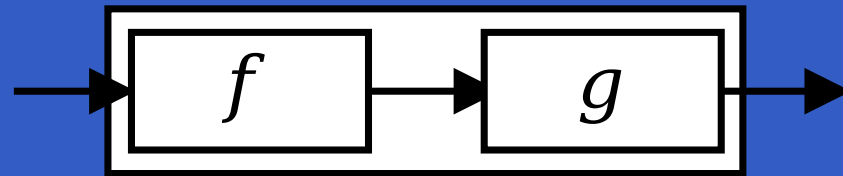
A **combinator** that captures this idea:

$$(\gg) :: SF\ a\ b \rightarrow SF\ b\ c \rightarrow SF\ a\ c$$

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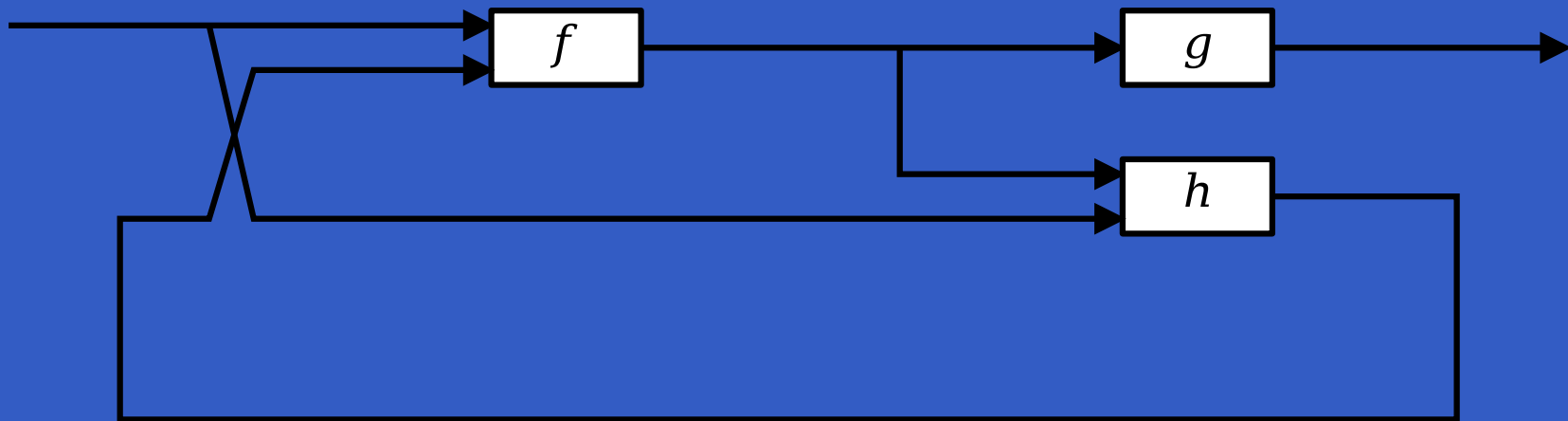
A **combinator** that captures this idea:

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Signal functions are the primary notion; signals a secondary one, only existing indirectly.

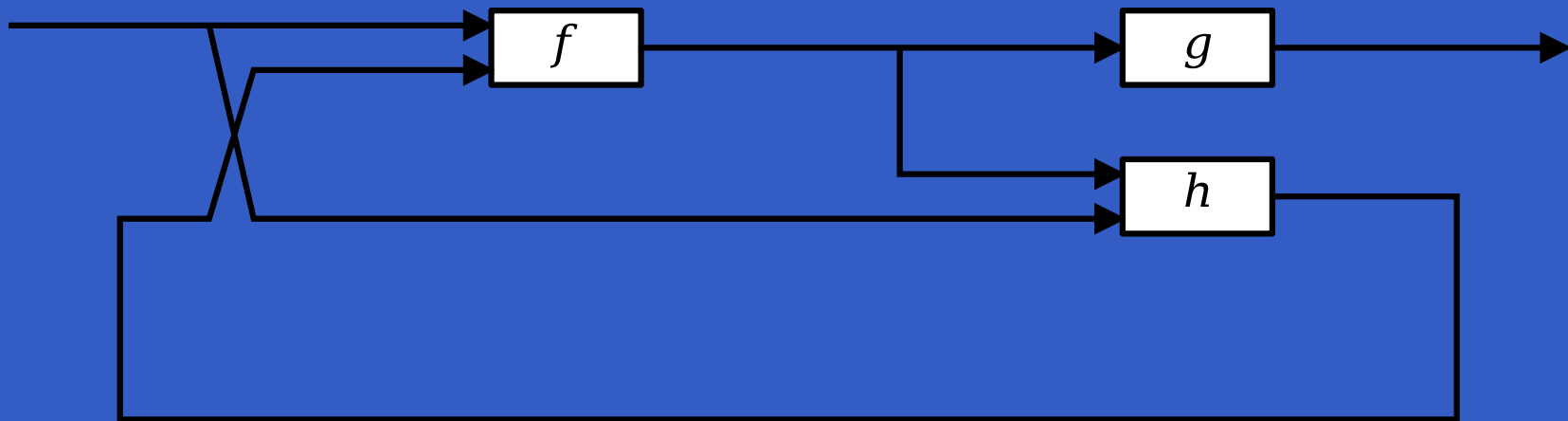
Systems

What about larger, more complicated networks?
How many combinators are needed?



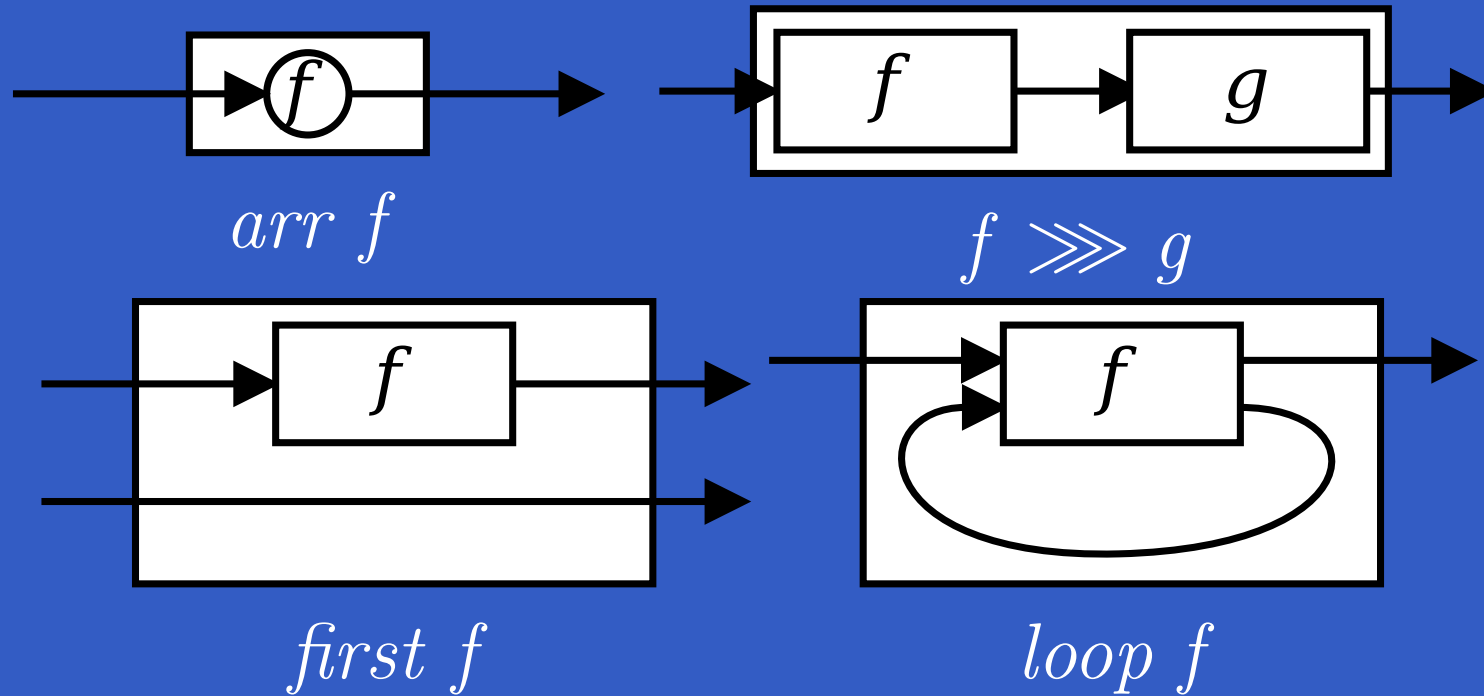
Systems

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John Hughes's **Arrow** framework provides a good answer!

The Arrow framework



$arr :: (a \rightarrow b) \rightarrow SF\ a\ b$

$(\ggg) :: SF\ a\ b \rightarrow SF\ b\ c \rightarrow SF\ a\ c$

$first :: SF\ a\ b \rightarrow SF\ (a, c)\ (b, c)$

$loop :: SF\ (a, c)\ (b, c) \rightarrow SF\ a\ b$

Oscillator from Pang-a-lambda

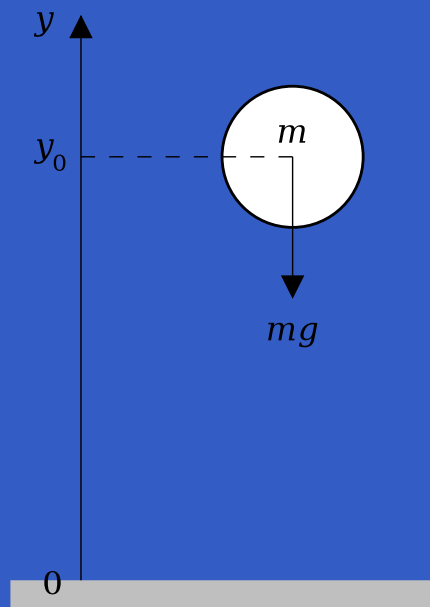
This oscillator determines the movement of blocks:

```
osci ampl period = proc _ → do  
  rec  
    let a =  $-(2.0 * \pi / \textit{period}) \uparrow 2 * p$   
    v ← integral ← a  
    p ←  $(\textit{ampl} +) \hat{\ll} \textit{integral} \leftarrow v$   
  return A ← p
```

Direct transliteration of standard equations.

A Bouncing Ball

Lots of bouncing balls in Pang-a-lambda!



$$y = y_0 + \int v \, dt$$

$$v = v_0 + \int -9.81$$

On impact:

$$v = -v(t-)$$

(fully elastic collision)

Modelling the Bouncing Ball: Part 1

Free-falling ball:

```
type Pos = Double
```

```
type Vel = Double
```

```
fallingBall :: Pos → Vel → SF () (Pos, Vel)
```

```
fallingBall y0 v0 = proc () → do
```

```
  v ← (v0+) ^<< integral ↯ - 9.81
```

```
  y ← (y0+) ^<< integral ↯ v
```

```
  returnA ↯ (y, v)
```

Events

Yampa models discrete-time signals by lifting the *co-domain* of signals using an option-type:

```
data Event a = NoEvent | Event a
```

Discrete-time signal = Signal (Event α).

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$$\text{data } \mathit{Event} \ a = \mathit{NoEvent} \mid \mathit{Event} \ a$$

Discrete-time signal = `Signal (Event a)`.

Some functions and event sources:

$$\mathit{tag} :: \mathit{Event} \ a \rightarrow b \rightarrow \mathit{Event} \ b$$
$$\mathit{after} :: \mathit{Time} \rightarrow b \rightarrow \mathit{SF} \ a \ (\mathit{Event} \ b)$$
$$\mathit{edge} :: \mathit{SF} \ \mathit{Bool} \ (\mathit{Event} \ ())$$

Modelling the Bouncing Ball: Part 2

Detecting when the ball goes through the floor:

fallingBall' ::

$Pos \rightarrow Vel \rightarrow SF () ((Pos, Vel), Event (Pos, Vel))$

fallingBall' *y0 v0* = **proc** () → **do**

yv@(*y*, -) ← *fallingBall* *y0 v0* ↦ ()

hit ← *edge* ↦ $y \leq 0$

returnA ↦ (*yv*, *hit* 'tag' *yv*)

Switching

Q: How and when do signal functions “start”?

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Switchers thus allow systems with **varying structure** to be described.

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Switchers thus allow systems with **varying structure** to be described.

Generalised switches allow composition of **collections** of signal functions. Can be used to model e.g. varying number of objects in a game.

The Basic Switch

Idea:

- Allows one signal function to be replaced by another.
- Switching takes place on the first occurrence of the switching event source.

switch::

$SF\ a\ (b,\ Event\ c)$

$\rightarrow (c \rightarrow SF\ a\ b)$

$\rightarrow SF\ a\ b$

The Basic Switch

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switch::

Initial SF with event source

$SF\ a\ (b,\ Event\ c)$

$\rightarrow (c \rightarrow SF\ a\ b)$

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The Basic Switch

Idea:

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switch::

Function yielding SF to switch into

$SF\ a\ (b,\ Event\ c)$

$\rightarrow (c \rightarrow SF\ a\ b)$

$\rightarrow SF\ a\ b$

Modelling the Bouncing Ball: Part 3

Making the ball bounce:

$bouncingBall :: Pos \rightarrow SF () (Pos, Vel)$

$bouncingBall\ y0 = bbAux\ y0\ 0.0$

where

$bbAux\ y0\ v0 =$

$switch\ (fallingBall'\ y0\ v0)\ \$\ \lambda(y, v) \rightarrow$

$bbAux\ y\ (-v)$

Game Objects

```
data Object = Object { objectName :: ObjectName
                      , objectKind  :: ObjectKind
                      , objectPos   :: Pos2D
                      , objectVel   :: Vel2D
                      ...
                      }
```

```
data ObjectKind = Ball ... | Player ... | ...
```

```
data ObjectInput = ObjectInput
  { userInput      :: Controller
  , collisions     :: Collisions
  }
```

Overall Game Structure

```
gamePlay :: [ListSF ObjectInput Object]  
           → SF Controller ([Object], Time)
```

```
gamePlay objs = loopPre [] $
```

```
  proc (input, cs) → do
```

```
    let oi = ObjectInput input cs
```

```
    ol ← dlSwitch objs ↯ oi
```

```
    let cs' = detectCollisions ol
```

```
    tLeft ← time ↯ ()
```

```
    returnA ↯ ((ol, tLeft), cs')
```

ListSF and *dlSwitch* are related abstractions that allow objects to die or spawn new ones.

And now for something different ...

Switched-on Yampa: Programming Modular Synthesizers in Haskell

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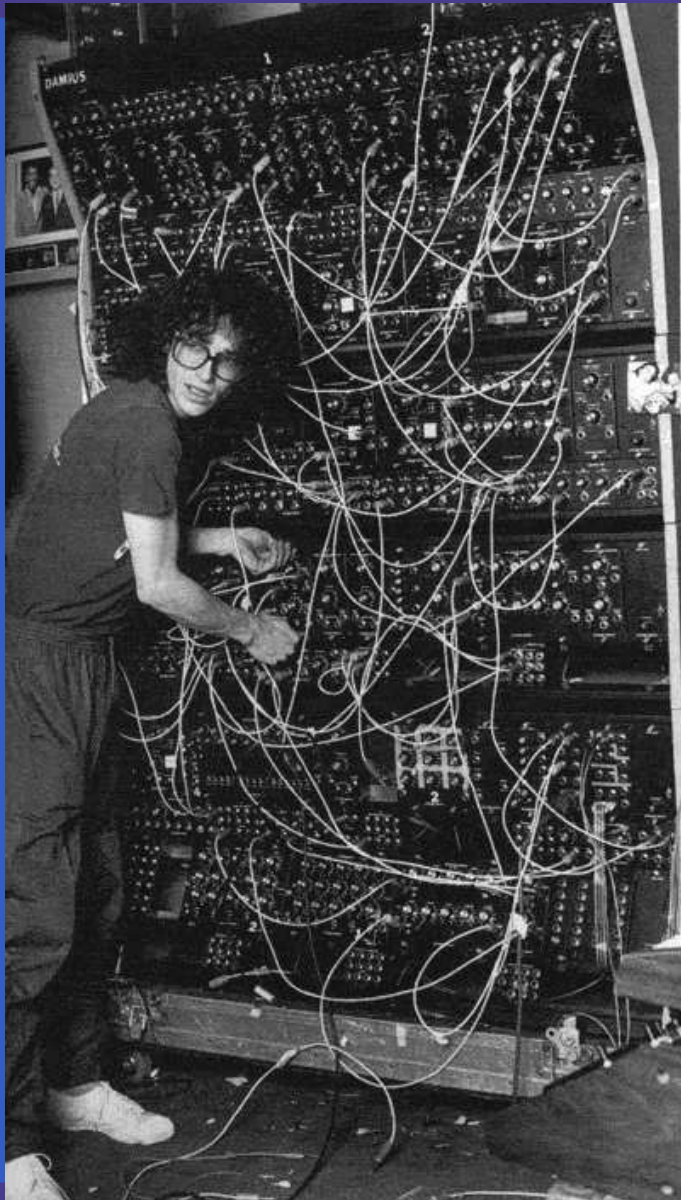
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Modular synthesizers?

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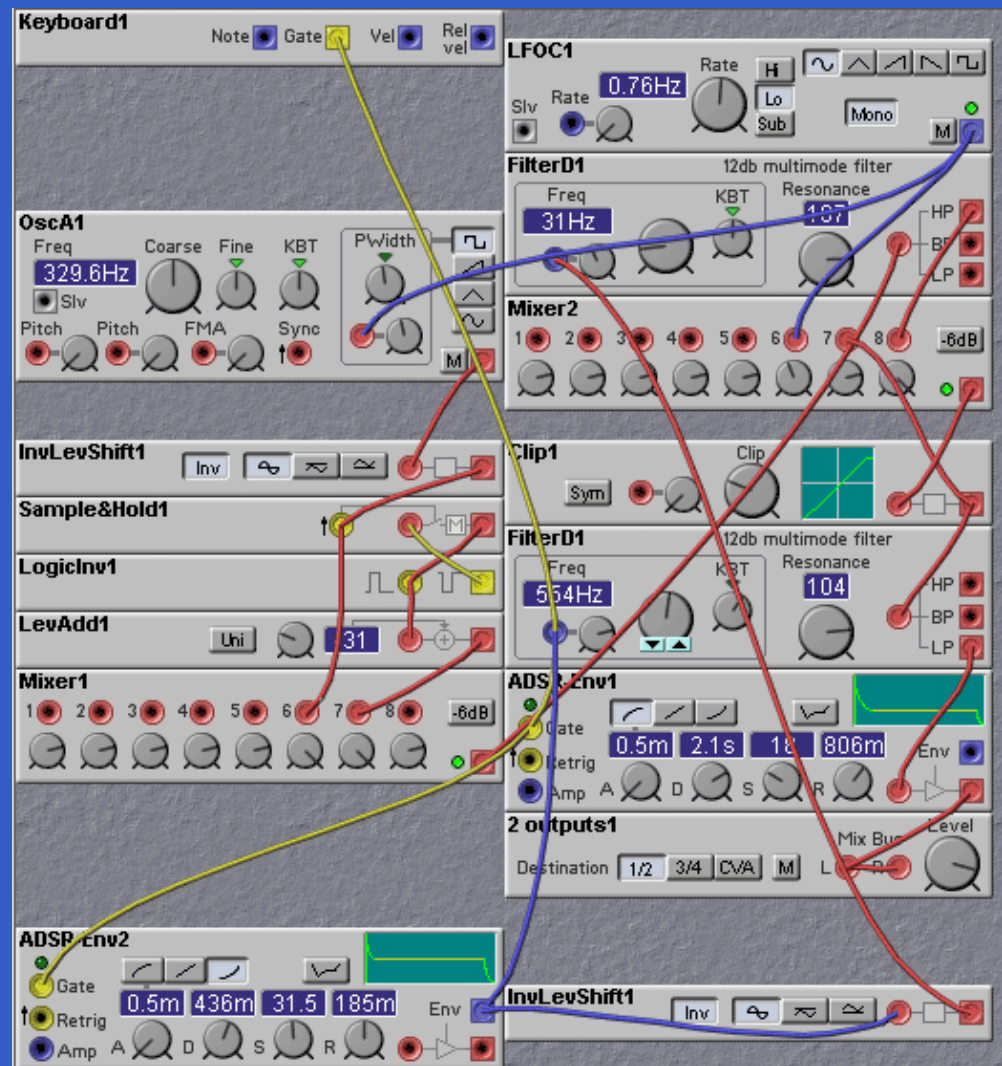


Modular synthesizers?



Steve Pocar, Toto,
with Polyfusion Syn-
thesizer

Modern Modular Synthesizers



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Where does Yampa enter the picture?



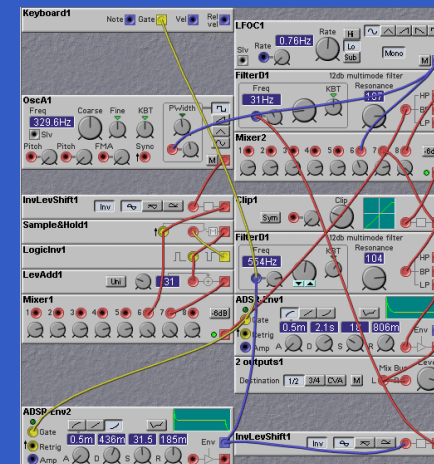
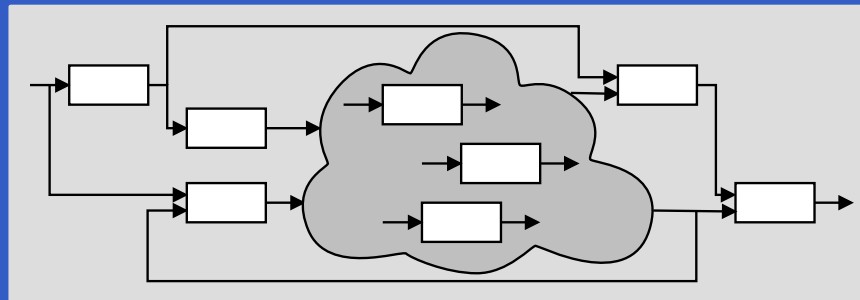
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Where does Yampa enter the picture?

- Music can be seen as a hybrid phenomenon. Thus interesting to explore a hybrid approach to programming music and musical applications.

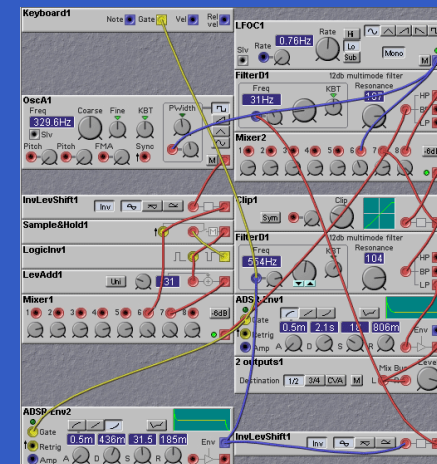
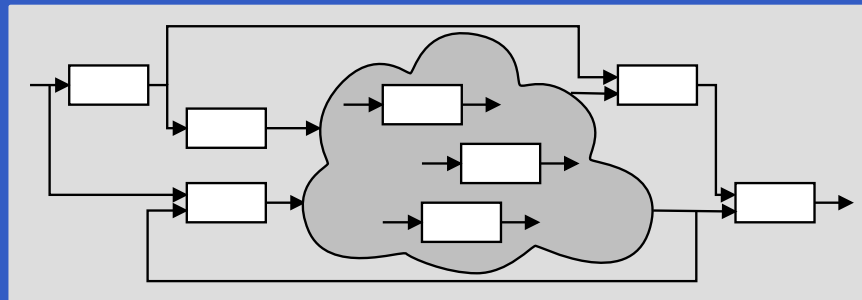
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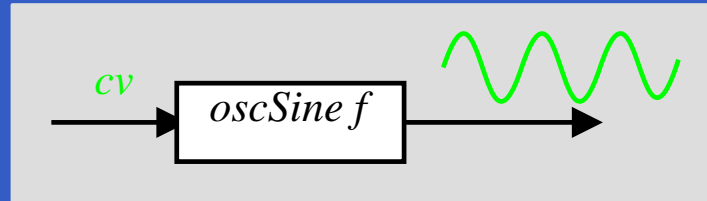
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- Fun application! Useful for teaching?

Example 1: Sine oscillator



oscSine :: Frequency \rightarrow SF CV Sample

oscSine f0 = **proc** cv \rightarrow do

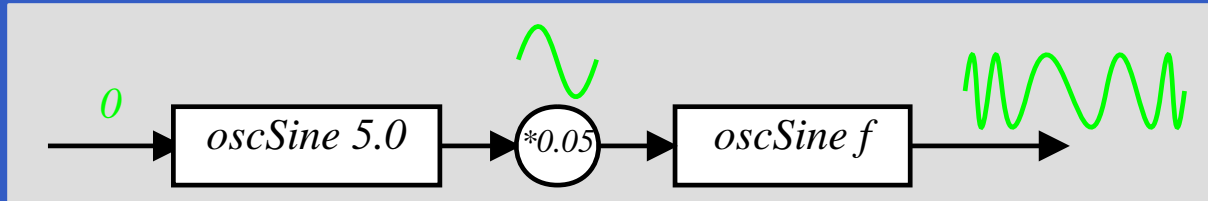
let $f = f0 * (2 ** cv)$

$\phi \leftarrow \text{integral} \leftarrow 2 * \pi * f$

return $A \leftarrow \sin \phi$

constant 0 \ggg *oscSine* 440

Example 2: Vibrato



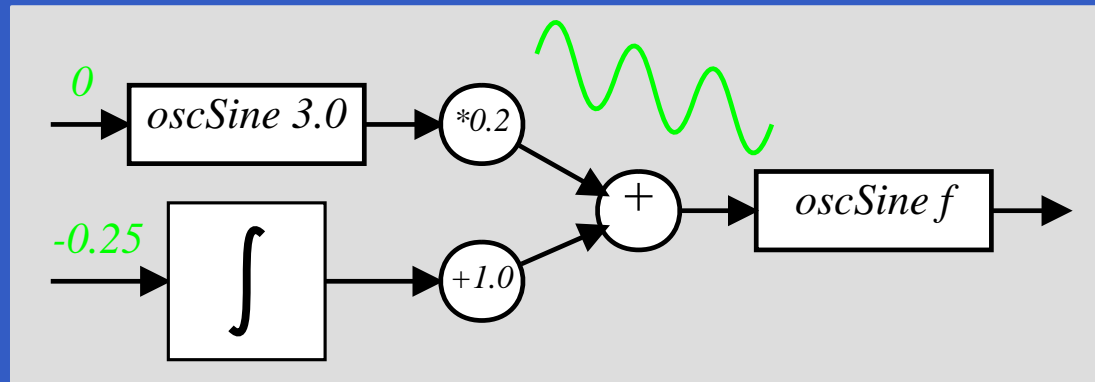
constant 0

≫≫ *oscSine 5.0*

≫≫ *arr (*0.05)*

≫≫ *oscSine 440*

Example 3: 50's Sci Fi



sciFi :: SF () Sample

sciFi = proc () → do

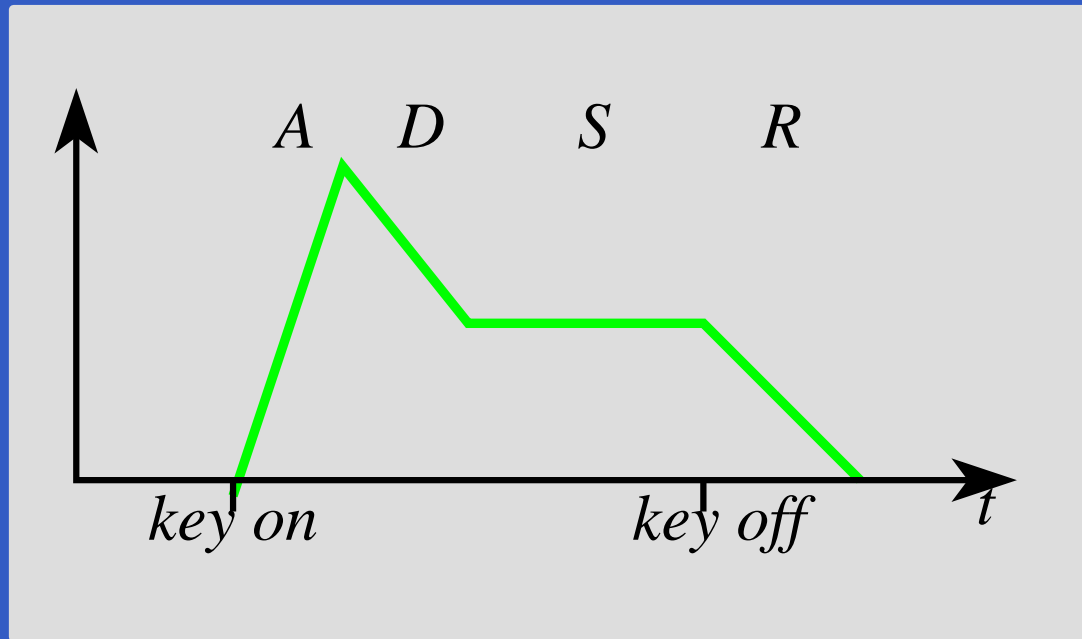
und ← arr (*0.2) <<< oscSine 3.0 — 0

swp ← arr (+1.0) <<< integral — -0.25

audio ← oscSine 440 — und + *swp*

return A — *audio*

Envelope Generators (1)



$envGen :: CV \rightarrow [(Time, CV)] \rightarrow (Maybe Int)$
 $\rightarrow SF (Event ()) (CV, Event ())$

$envEx = envGen 0 [(0.5, 1), (0.5, 0.5), (1.0, 0.5), (0.7, 0)]$
(Just 3)

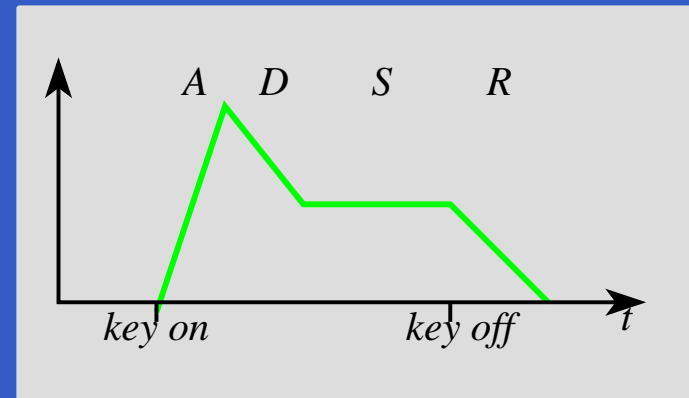
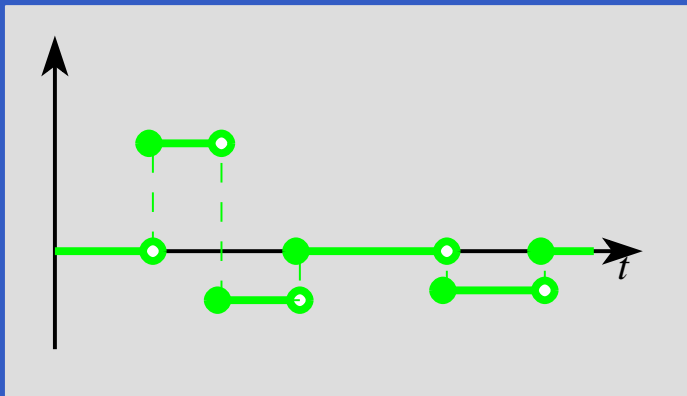
Envelope Generators (2)

How to implement?

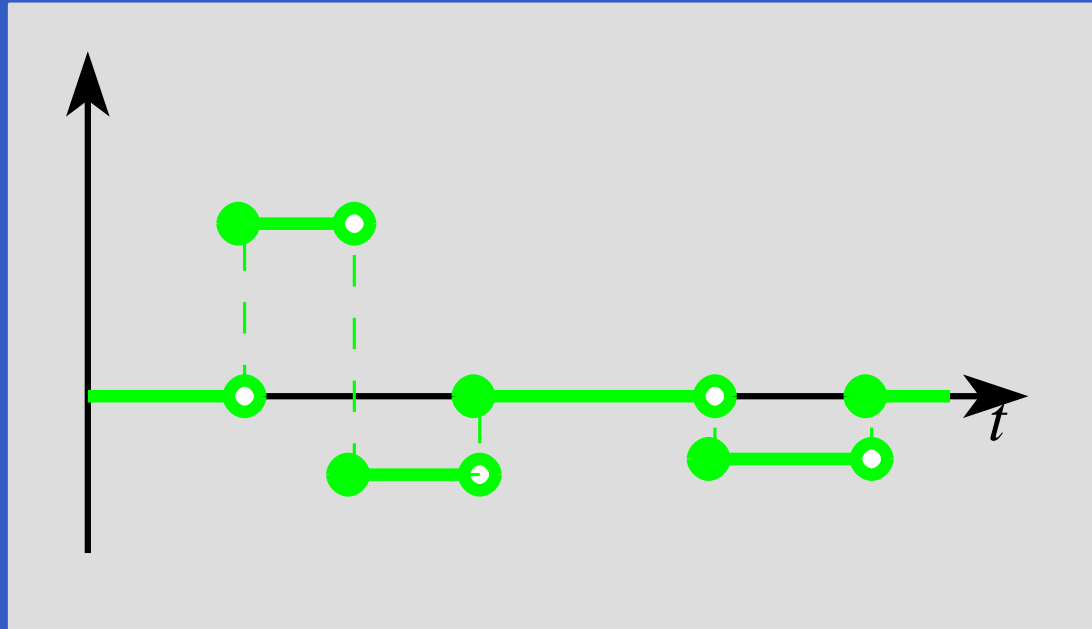
Envelope Generators (2)

How to implement?

Integration of a step function yields suitable shapes:



Envelope Generators (3)



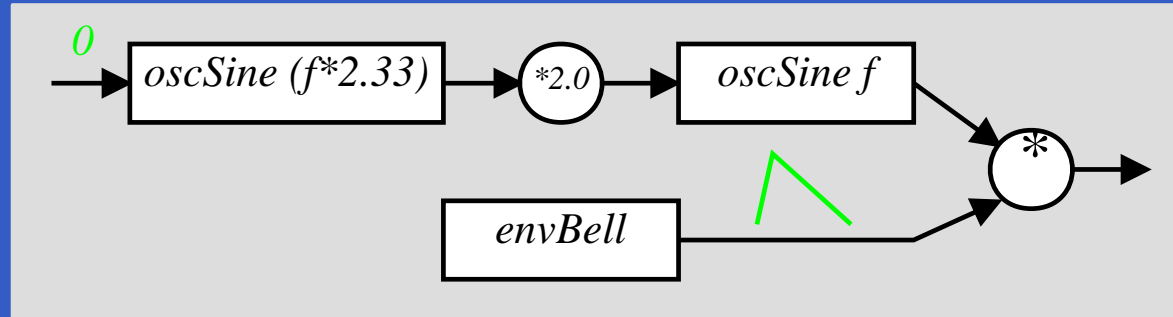
$afterEach :: [(Time, b)] \rightarrow SF\ a\ (Event\ b)$

$hold \quad \quad :: a \rightarrow SF\ (Event\ a)\ a$

$steps = afterEach [(0.7, 2), (0.5, -1), (0.5, 0), (1, -0.7), (0.7, 0)]$

$\ggg hold\ 0$

Example 4: Bell



$bell :: \text{Frequency} \rightarrow SF () (\text{Sample}, \text{Event})$

$bell f = \mathbf{proc} () \rightarrow \mathbf{do}$

$m \leftarrow oscSine (2.33 * f) \prec 0$

$audio \leftarrow oscSine f \prec 2.0 * m$

$(ampl, end) \leftarrow envBell \prec noEvent$

$return A \prec (audio * ampl, end)$

Example 5: Tinkling Bell

tinkle :: SF () Sample

tinkle = (repeatedly 0.25 84

 >>> constant ())

 &&arr (fmap (bell ◦ midiNoteToFreq))

 >>> rSwitch (constant 0))

Example 6: Playing simultaneous notes

```
mysterySong :: SF () (Sample, Event ())  
mysterySong = proc _ → do  
  t ← tinkle ↯ ()  
  m ← mystery ↯ ()  
  return A ↯ (0.4 * t + 0.6 * m)
```

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Switched-on Yampa?

Switched-on Yampa?



Software and paper: www.cs.nott.ac.uk/~ggg

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Conclusions

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Conclusions

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- It allows systems to be described in terms of whole values varying over time.
- Not covered in this talk:
 - Not everything fit easily into the FRP paradigm: e.g., interfacing to existing GUI toolkits, other imperative APIs.
 - But also such APIs can be given a “whole-value treatment” to improve the fit within a declarative setting. E.g. ***Reactive Values and Relations***.