

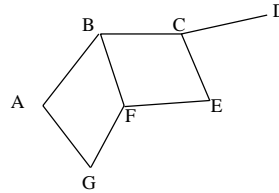
## Graphs

Plan of the lecture:

- What is a graph
- What are they used for
- Graph problems
- Two ways of implementing graphs

## Definition of a graph

A graph is a set of *nodes*, or *vertices*, connected by *edges*.



## Applications of Graphs

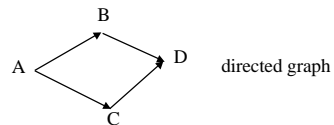
Graphs can be used to represent

- networks
- flow charts
- states of automaton / program
- tasks in some project (some of which should be completed before others), so edges correspond to prerequisites.

## Directed and Undirected Graphs

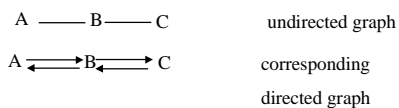
Graphs can be

- undirected (edges don't have direction).
- directed (edges have direction).



## Directed and Undirected Graphs

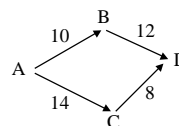
Undirected graphs can be represented as directed graphs where for each edge  $(X,Y)$  there is a corresponding edge  $(Y,X)$ .



## Weighted and Unweighted Graphs

Graphs can also be

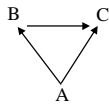
- unweighted (as in the previous examples)
- weighted (edges have weights).



## Notation

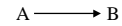
- Set  $V$  of *vertices* (nodes)
- Set  $E$  of *edges* ( $E \subseteq V \times V$ )

Example:  $V = \{A, B, C\}$ ,  $E = \{(A,B), (A,C), (B,C)\}$ :



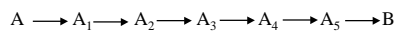
## Adjacency relation

- Node  $B$  is *adjacent* to  $A$  if there is an edge from  $A$  to  $B$ .



## Paths and reachability

- A *path* from  $A$  to  $B$  is a sequence of vertices  $A_1, \dots, A_n$  such that there is an edge from  $A$  to  $A_1$ , from  $A_1$  to  $A_2$ , ..., from  $A_n$  to  $B$ .



- A vertex  $B$  is *reachable* from  $A$  if there is a path from  $A$  to  $B$

## More Terminology

- A *cycle* is a path from  $u$  to itself
- Graph is *acyclic* if it does not have cycles
- Graph is *connected* if there is a path between every pair of vertices
- Graph is *strongly connected* if there is a path in both directions between every pair of vertices

## Some graph problems

- searching a graph for a vertex
- finding a path in the graph (from one vertex to another); finding the shortest path between two vertices
- cycle detection

## More graph problems

- topological sort (finding a linear sequence of vertices which agrees with the direction of edges in the graph; applied for scheduling tasks in a project)
- minimum spanning tree (deleting as many edges in a graph as possible, so that all vertices are still connected by shortest possible edges. Useful in circuit design.)

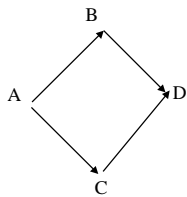
## How to implement a graph

- As with lists, two approaches:
  - using a static indexed data structure
  - using a dynamic data structure

## Static implementation:Adjacency Matrix

- Store nodes in the array: each node is associated with an integer (array index).
- Represent information about the edges using a two dimensional array, where  
 $\text{array}[i][j] == 1$   
iff there is an edge from node  $i$  to node  $j$ .

## Example



A	B	C	D
0	1	2	3

node indices

	0	1	2	3
0	0	1	1	0
1	0	0	0	1
2	0	0	0	1
3	0	0	0	0

adjacency matrix

## Weighted graphs

- For weighted graphs, place weights in matrix (for no edge, -1 or Integer.MAX\_VALUE)

## Example implementation

- If we know which nodes the graph is going to have, and the number of nodes is not going to change, we can implement a graph as follows.
- **ArrayGraph** class has an array to keep nodes (Objects) and a two dimensional array to keep 0s and 1s depending on whether there is an edge between corresponding nodes.

## ArrayGraph class

```
public class ArrayGraph {
    int[][] matrix; // adjacency matrix
    Object[] positions;
    public ArrayGraph(Object[] nodes) {
        positions = nodes;
        matrix = new
            int[nodes.length][nodes.length];
    } // this fills matrix with 0s
```

## ArrayGraph class

```
public void addEdge(Object x, Object y) {
    if (indexOf(x) == -1 | indexOf(y) == -1)
        return;
    else matrix[indexOf(x)][indexOf(y)] = 1;
}
/* This method does nothing if nodes do
not exist. indexOf method follows. */
```

## ArrayGraph class

```
// to find o's position in the matrix...
public int indexOf(Object o) {
    for(int i=0; i<positions.length; i++){
        if(positions[i].equals(o)) {
            return i;
        }
    }
    return -1; // o is not in the graph
}
```

## Notes on ArrayGraph class

- Keeping edge labels in an array is a bit awkward (have to look through the array each time to find out node's position in the matrix)
- Alternative solution could be using a HashMap: the Object (node) is the key, and Integer object holding its index is the value.
- When we need to know the node o's position we could do

```
int pos =
((Integer) map.get(o)).intValue();
```

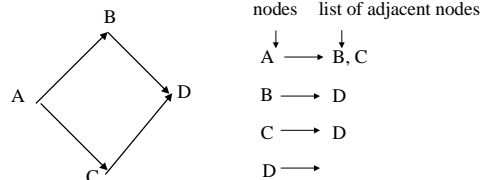
## Disadvantages of adjacency matrices

- Sparse graphs: few edges for number of vertices.
- Leads to many zero entries in adjacency matrix: wastes space, makes many algorithms less efficient (to find nodes adjacent to a given node, have to iterate through the whole row, even if there are few 1s there).
- Also, if the number of nodes in the graph may change, matrix representation is too inflexible (especially if we don't know the maximal size of the graph).

## Adjacency List

- For every vertex, keep a list of adjacent vertices.
- Represent a graph as a list or array of such lists: this is called an adjacency list implementation.

## Adjacency list



## Adjacency List Implementation

- Graph can be implemented as a list of lists, or any other structure holding vertices and lists of their neighbours.
- The simplest example is the **LinkedListGraph** class: it keeps a list of **GraphNode** objects. Each **GraphNode** object keeps an **Object** and a neighbours list (of **GraphNode** objects).

## GraphNode class

```
class GraphNode {
    Object label;
    LinkedList neighbours;

    GraphNode(Object o) {
        label = o;
        neighbours = new LinkedList();
    }
}
```

## LinkedListGraph class

```
import java.util.*;
public class LinkedListGraph {
    LinkedList nodes;
    public LinkedListGraph() {}
    // constructor sets fields to null
    public void addNode(Object o) {
        nodes.add(new GraphNode(o));
    }
}
```

## LinkedListGraph class

```
public boolean containsNode(Object o) {
    ListIterator li = nodes.listIterator();
    while (li.hasNext()) {
        GraphNode n = (GraphNode) li.next();
        if (n.label.equals(o)) {
            return true;
        }
    }
    return false;
}
```

## LinkedListGraph class

- ... other methods you would have to write yourself, but the idea should be clear:
- given an object (or two objects), find their corresponding **GraphNode** objects in the list, and then modify their neighbours lists etc.
- To implement graph traversal algorithms, we may need to add extra fields to the **GraphNode** class, for example a boolean flag to say that we have seen this node before.

## Notes on LinkedListGraph class

- Again, iterating through the list of nodes every time we need to find a **GraphNode** object corresponding to an **Object o** is awkward.
- One option (there are many others) is to use a **HashMap**: **Object o** is the key, and its neighbours list (of **Objects**) is the value. This way, don't need a **GraphNode** class.

## HashMap implementation

Key	Value
A	Linked list containing B,C
B	Linked list containing D
C	Linked list containing D
D	Empty linked list

## Summary and Reading

- Graphs can be used in many applications (anywhere where diagrams and maps are used).
- Graphs can be implemented with adjacency matrices or adjacency lists.
- For the formal coursework, have a look at Java Collections API to choose a suitable data structure for your implementation.
- Have a look at Shaffer, chapter 7 for graph terminology. His implementation of graphs assumes fixed size graphs (storing ints).