Homotopy Type Theory

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Functional Programming Laboratory Away Day

8th July 2011

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What is it all about?

A connection...





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Type Theory

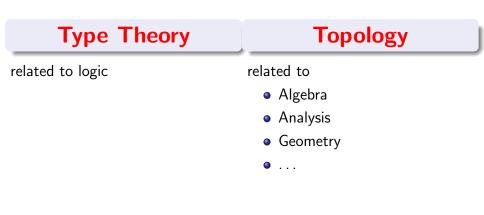
Topology

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related to logic

What is it all about?

A connection...



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Sets often have a natural notion of *open* subsets, e.g. in \mathbb{R} :

$$(1,2) := \{x \mid 1 < x < 2\}$$
 is open, but $[1,2] := \{x \mid 1 \le x \le 2\}$ is not.

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Sets often have a natural notion of *open* subsets, e.g. in \mathbb{R} :

 $(1,2) := \{x \mid 1 < x < 2\}$ is open, but $[1,2] := \{x \mid 1 \le x \le 2\}$ is not.

Definition (Continuity of $f : X \rightarrow Y$):

f is continuous iff inverse images of open sets are open, i.e. if $V \subset Y$ is open, so is $f^{-1}(V)$.

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• functions $\mathbb{R} \to \mathbb{R}$. "Definition": $\forall x, \epsilon > 0. \exists \delta > 0. |x - x_0| < \delta \Rightarrow |f(x) - f(x_0)| < \epsilon$

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all notions can be broken down to:

Inverse images of open sets are open.

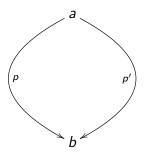
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$$\frac{a, b: A}{a \equiv b \quad \text{Type}}$$

 $p : a \equiv b$ $p^{-1} : b \equiv a$ $q : b \equiv c$ $q \circ p : a \equiv c$

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 $\operatorname{refl}_a : a \equiv a$



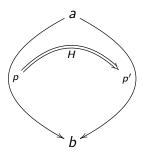
for example:

• a := b := x

•
$$p := p' := \operatorname{refl}_x$$

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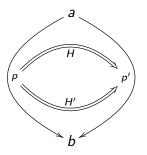


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for example:

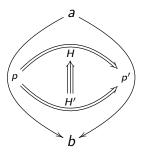
• *a* := *b* := *x*

•
$$p := p' := \operatorname{refl}_x$$

•
$$H := H' := \operatorname{refl}_{\operatorname{refl}_x}$$

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for example:

- *a* := *b* := *x*
- $p := p' := \operatorname{refl}_x$
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• refl_{refl}refl_x

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Structures:

Topological spaces

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Topological spaces

 \rightarrow Special case: Hausdorff spaces (or T_2)

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Structures:

Topological spaces

- \rightarrow Special case: Hausdorff spaces (or T_2)
- $\rightarrow \rightarrow$ Special case of this special case: Metric spaces

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- $\rightarrow \rightarrow \rightarrow$ Even much more special: Normed vector spaces

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 $\rightarrow \rightarrow \rightarrow \rightarrow$...and finally: $\mathbb{R}^n,$ or just subsets of it!

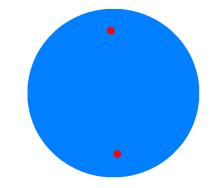


a topological space - we call it X

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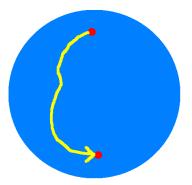
two terms



two points

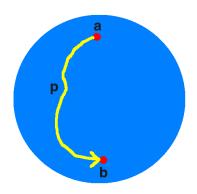
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a path

a, b : X $p : a \equiv b$

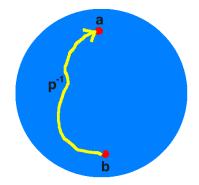


$$a, b \in X$$

 $p : [0, 1] \rightarrow X$
 $p(0) = a$
 $p(1) = b$

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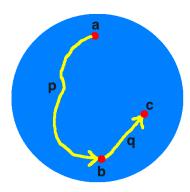
 p^{-1} : $b \equiv a$



$$p^{-1}: [0,1] o X \ p^{-1}(t) = p(1-t)$$

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 $p: a \equiv b$ $q: b \equiv c$



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$$p : [0,1] \rightarrow X$$

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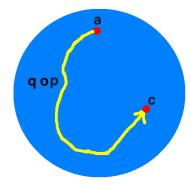
$$q : [0,1] \rightarrow X$$

$$q(0) = b$$

$$q(1) = c$$

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$$egin{aligned} q \circ p &: \ [0,1] & o X \ x \mapsto \ \left\{ egin{aligned} p(2x), x < 0.5 \ q(2x-1), ext{else} \end{aligned}
ight. \end{aligned}$$

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 $q \circ p : a \equiv c$

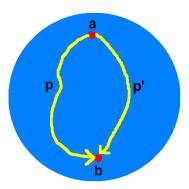
Another set

 $a \equiv c \text{ not}$ inhabited

not path-connected

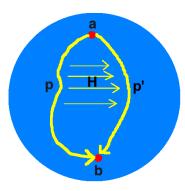
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$$p, p' : a \equiv b$$



$p,p'\,:\,[0,1]\to X$

$$H : p \equiv p'$$



$$H : [0,1]^2 \rightarrow X$$
$$H(0,\cdot) = p$$
$$H(1,\cdot) = p'$$
$$H(t,0) = a$$
$$H(t,1) = b$$

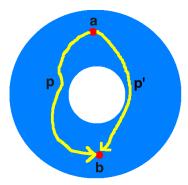
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 $p : [0,1]^1 \rightarrow X$ $a : [0,1]^0 \rightarrow X$

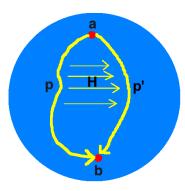
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$$H : p \equiv p'$$

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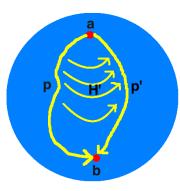


$$H : p \equiv p'$$



$$H : [0,1]^2 \rightarrow X$$
$$H(0,\cdot) = p$$
$$H(1,\cdot) = p'$$
$$H(t,0) = a$$
$$H(t,1) = b$$

$$H' \cdot p = p'$$

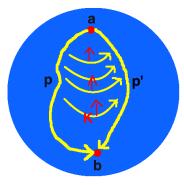


 $H' : [0,1]^2 \to X$ $H'(0, \cdot) = p$ $H'(1, \cdot) = p'$ H'(t,0) = aH'(t,1) = b

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$egin{array}{ll} \mathcal{K} \ \colon \ [0,1]^3 o X \ \mathcal{K}(0,\cdot,\cdot) = \mathcal{H}' \end{array}$

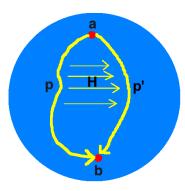
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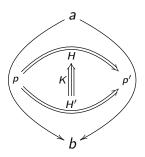
$K : H' \equiv H$

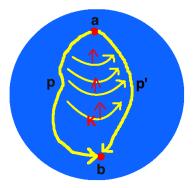
$$H : p \equiv p'$$



$$H : [0,1]^2 \rightarrow X$$
$$H(0,\cdot) = p$$
$$H(1,\cdot) = p'$$
$$H(t,0) = a$$
$$H(t,1) = b$$

Putting it together





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So, which types can we get?



any CW complex?

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Where is it going?

All has been done in abstract homotopy theory.



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All has been done in abstract homotopy theory.

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What I (at the moment) hope:

- Creating a simple model
- that is complete
- and easy to understand and to use

Where is it going?

- e.g. for this problem (Thorsten):
 - subst-refl P (subst P (refl x) p) and
 - cong (subst P (refl x)) (subst-refl P p) both prove
 - subst P (refl x) (subst P (refl x) p) \equiv subst P (refl x) p But are they equal?

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