

An Introduction to Type Theory

Practical 1

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Thorsten Altenkirch
University of Nottingham

tutch - example of a proof checker

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tutch *tutorial proof checker*

is a proof checker for educational purposes implemented by
Andreas Abel (Munich)
when he was working for Frank Pfenning at Carnegie
Mellon University, USA

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tutch syntax

| Logic | tutch |
|-------------------|-------|
| \wedge | & |
| \vee | |
| \neg | ~ |
| \Rightarrow | => |
| T | T |
| F | F |
| \Leftrightarrow | <=> |

$$A \wedge B \Rightarrow B \wedge A$$

```
proof comAnd : A & B => B & A =
begin
[A & B;
 A;
 B;
 B & A];
A & B => B & A
end;
```

$$A \wedge B \Rightarrow B \wedge A$$

```
proof comAnd : A ∧ B ⇒ B ∧ A =
begin
[A ∧ B;
 A;
 B;
 B ∧ A];
A ∧ B ⇒ B ∧ A
end;
```

Running tutch

1. Create a text file comAnd.tut
2. Run tutch:

```
[txa@jacob mytutch]$ tutch comAnd.tut
TUTCH 0.51 beta, $Date: 2000/10/27 17:08:48 $
[Opening file comAnd.tut]
Proving Cla: A ∧ B ⇒ B ∧ A ...
QED
[Closing file comAnd.tut]
```

tutch verbose

```
[txa@jacob mytutch]$ tutch -v comAnd.tut
TUTCH 0.51 beta, $Date: 2000/10/27 17:08:48 $
[Opening file comAnd.tut]
```

Proving comAnd: $A \wedge B \Rightarrow B \wedge A \dots$

```
1 [ A ∧ B;
2   A;                                by AndEL 1
3   B;                                by Ander 1
4   B ∧ A ];                            by AndI 3 2
5 A ∧ B ⇒ B ∧ A                      by ImpI 4
```

QED

An incomplete proof

```
proof comAnd : A ∧ B ⇒ B ∧ A =
begin
[A ∧ B;
 B ∧ A];
A ∧ B ⇒ B ∧ A
end;
```

tutch's error messages

```
[txa@jacob mytutch]$ tutch comAnd.tut
TUTCH 0.51 beta, $Date: 2000/10/27 17:08:48 $  

[Opening file comAnd.tut]
Proving comAnd: A ∧ B ⇒ B ∧ A ...
comAnd.tut:4.2-4.7 Error:
Unjustified line A ∧ B   ⊢   B ∧ A
Assuming this line, checking remainder...
Proof incomplete
[Closing file comAnd.tut]
```

Rules for \wedge

AndI

A;
B;
A \wedge B;

AndEL

A \wedge B;
A;

AndER

A \wedge B;
B;

Rules for \Rightarrow

Impl

```
[ A;  
  :  
  B ] ;  
A  $\Rightarrow$  B ;
```

ImpE

```
A  $\Rightarrow$  B ;  
A ;  
B ;
```

The simplest proof?

```
proof I: A ⇒ A =  
begin  
[ A;  
  A ];  
A ⇒ A  
end;
```

tutch output

Proving I: A \Rightarrow A . . .

1 [A;

2 A] ; by Hyp 1

3 A \Rightarrow A by ImpI 2

QED

Rules for \vee

OrI_L

$$\begin{array}{c} A; \\ A \vee B; \end{array}$$

OrI_R

$$\begin{array}{c} B; \\ A \vee B; \end{array}$$

Rules for \vee (cont)

OrE

```
[ A;  
  :  
  C ] ;  
[ B;  
  :  
  C ] ;  
A  $\vee$  B;  
C;
```

$$A \vee B \Rightarrow B \vee A$$

```
proof comOr : A ∨ B ⇒ B ∨ A =
begin
[ A ∨ B;
  [ A;
    B ∨ A];
  [ B;
    B ∨ A];
  B ∨ A];
A ∨ B ⇒ B ∨ A
end;
```

tutch output

Proving comOr: $A \vee B \Rightarrow B \vee A \dots$

- 1 [A \vee B;
- 2 [A;
- 3 B \vee A] ; by OrIR 2
- 4 [B;
- 5 B \vee A] ; by OrIL 4
- 6 B \vee A] ; by OrE 1 3 5
- 7 A \vee B \Rightarrow B \vee A by Impl 6

QED

T and F

TrueL

T ;

FalseE

F ;

A ;

\neg and \Leftrightarrow

$\neg A$ is defined as $A \Rightarrow \text{F}$

$A \Leftrightarrow B$ is defined as $(A \Rightarrow B) \wedge (B \Rightarrow A)$

$$\neg(A \wedge \neg A)$$

```
proof incons :  $\neg(A \wedge \neg A) =$ 
begin
[ A  $\wedge$   $\neg$  A;
  A;
   $\neg$  A;
  F ] ;
 $\neg(A \wedge \neg A)$ 
end;
```

Propositional logic

Prove the following in intuitionistic propositional logic using tutch (prop.req):

1. $A \Rightarrow A \wedge A$
2. $A \vee A \Rightarrow A$
3. $(A \Rightarrow B \Rightarrow C) \Rightarrow (A \Rightarrow B) \Rightarrow (A \Rightarrow C)$
4. $A \wedge (B \vee C) \Leftrightarrow (A \wedge B) \vee (A \wedge C)$
5. $\neg A \wedge \neg B \Leftrightarrow \neg(A \vee B)$

Using prop.req

You can check your proofs by running tutch under unix by typing

```
tutch -r ./prop.req -v prop.tut
```

in the shell window and edit your proofs until you get the message

Congratulations! All problems solved!

Predicate logic: tutch syntax

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Tutch is *typed*, we write

$$\forall x : T.P(x)$$

$$\exists x : T.P(x)$$

where T is a type (e.g. nat).

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$$\forall x : T.P(x)$$

$$\exists x : T.P(x)$$

where T is a type (e.g. nat).

In the moment we just use a type variable (e.g. t).

Rules for \forall

ForallI

$$\begin{array}{l} [x : t ; \\ \vdots \\ P] ; \\ \forall x : t . P ; \end{array}$$

ForallE

$$\begin{array}{l} \forall x : t . P ; \\ u : t ; \\ P [x := u] \end{array}$$

Example of a proof using \forall

```
proof allCom :  
  ( $\forall x:t. \forall y:t. P(x,y) \Rightarrow (\forall y:t. \forall x:t. P(x,y))$ ) =  
begin  
  [ $\forall x:t . \forall y:t. P(x,y)$ ];  
  [ $y:t$ ;  
   [ $x:t$ ;  
    [ $\forall y:t. P(x,y)$ ;  
      $P(x,y)$ ]];  
   [ $\forall x:t. P(x,y)$ ]];  
  [ $\forall y:t. \forall x:t. P(x,y)$ ];  
  ( $\forall x:t. \forall y:t. P(x,y) \Rightarrow (\forall y:t. \forall x:t. P(x,y))$ );  
end;
```

Predicate logic

Prove the following in intuitionistic propositional logic using tutch (pred.req)

1. $(\forall x : t.P(x)) \Rightarrow (\forall y : t.P(y))$
2. $(\forall x : t.P(x) \wedge Q(x)) \Leftrightarrow (\forall x : t.P(x)) \wedge (\forall x : t.Q(x))$
3. $(\exists x : t.P(x) \wedge Q) \Rightarrow (\exists x : t.P(x)) \wedge Q$
4. $(\exists x : t.P(x) \vee Q(x)) \Leftrightarrow (\exists x : t.P(x)) \vee (\exists x : t.Q(x))$