

λ -calculus

Other Data Structures

Boolean operations can be defined in terms of **if**

Booleans

$$\begin{aligned} \text{true} &:= \lambda x. \lambda y. x \} \text{ same as} \\ \text{false} &:= \lambda x. \lambda y. y \end{aligned} \quad \text{Projections}$$

Pairs

The if-then-else operation can be implemented as

$$\text{if } := \lambda b. \lambda u. \lambda v. buv$$

If t_1 and t_2 are λ -terms we can represent the pair of them by

$$\langle t_1, t_2 \rangle := \lambda x. x t_1 t_2$$

For example

$$\text{or} := \lambda a. \lambda b. \text{if } a \text{ true } b$$

Projections are just applications

to Booleans:

$$\begin{aligned} \text{if true } t_1 t_2 &\sim^* t_1 \\ \text{if false } t_1 t_2 &\sim^* t_2 \end{aligned}$$

Vectors can be repeated pairs:

$$\begin{aligned}\text{first} &:= \lambda p. p \text{ true} \\ \text{second} &:= \lambda p. p \text{ false}\end{aligned}$$

Check that they work:

$$\text{first } \langle t_1, t_2 \rangle$$

$$= (\lambda p. p \text{ true}) \langle t_1, t_2 \rangle$$

$$\rightsquigarrow \langle t_1, t_2 \rangle \text{ true}$$

$$= (\lambda x. x t_1 t_2) (\lambda x. \lambda y. x)$$

$$\rightsquigarrow (\lambda x. \lambda y. x) t_1 t_2$$

$$\rightsquigarrow^* t_1$$

Similarly

$$\text{second } \langle t_1, t_2 \rangle \rightsquigarrow^* t_2$$

or directly defined similarly:

$$\langle t_1, t_2, t_3 \rangle := \lambda x. x t_1 t_2 t_3$$

Lists

We may think of defining also lists as repeated pairs:

$$[t_1, t_2, \dots, t_n] := \langle t_1, \langle t_2, \dots \rangle \rangle$$

But how about the empty list?

We could give it a conventional

value: nil := false

What about the test: isnil ?

A more convenient way

in the spirit of Church numerals:

$$[t_1, t_2, t_3] := \lambda f. \lambda x. f t_1 (f t_2 (f t_3 x))$$

Verify that it is correct:

$$\text{isnil nil} \rightsquigarrow^* \text{true}$$

$$\text{isnil (cons h t)} \rightsquigarrow^* \text{false}$$

In general:

$$\text{nil} := \lambda f. \lambda x. x$$

$$\text{cons} := \lambda h. \lambda t.$$

$$\lambda f. \lambda x. f h (t f x)$$

Exercise:

Check what

$$\text{cons } t_1 (\text{cons } t_2 \text{ nil})$$

reduces to

The empty list test:

$$\text{isnil} := \lambda e. e (\lambda y. \lambda z. \text{false}) \text{ true}$$

Exercise:

Write head and tail functions
extracting the first element
and the rest of a non-empty
list.

(Doesn't matter what it does
on nil.)

Exercise:

- Define a function `sum` that adds up all elements of a list of numerals.
- Define a function `mirror` that reverses a list.

Can we do infinite lists?

An infinite list (or stream) is a neverending sequence:

$$a_0 \Delta a_1 \Delta a_2 \Delta \dots$$

Idea: represent it as

$$\lambda f. f a_0 (f a_1 (f a_2 \dots))$$

But it can't actually be infinitely long.

The stream $0 \Delta 1 \Delta 2 \Delta \dots$ should be a term t_0 such that:

$$t_0 \rightsquigarrow^* \lambda f. f 0 t_1$$

$$\rightsquigarrow^* \lambda f. f 0 (f \mid t_2)$$

$$\rightsquigarrow^* \lambda f. f 0 (f \mid (f 2 t_3)) \dots$$