

## Recap: Formal Definition of DFA

Formally, a **Deterministic Finite Automaton** or **DFA** is defined by a 5-tuple

$$(Q, \Sigma, \delta, q_0, F)$$

where

- $Q$  : **Finite** set of States
- $\Sigma$  : Alphabet (finite set of symbols)
- $\delta \in Q \times \Sigma \rightarrow Q$  : Transition Function
- $q_0 \in Q$  : Initial or Start State
- $F \subseteq Q$  : Accepting (or Final) States

## Recap: Language of a DFA

The **language**  $L(A)$  defined by a DFA  $A$  is the set or words **accepted** by the DFA. For a DFA

$$A = (Q, \Sigma, \delta, q_0, F)$$

the language is defined by

$$L(A) = \{ w \in \Sigma^* \mid \hat{\delta}(q_0, w) \in F \}$$

# G52MAL

## Machines and Their Languages

### Lecture 3

### *Non-deterministic Finite Automata (NFA)*

Henrik Nilsson

University of Nottingham

## Recap: Extended Transition Function

The **Extended Transition Function** is defined on a state and a **word** (string of symbols) instead of on a single symbol.

For a DFA  $A = (Q, \Sigma, \delta, q_0, F)$ , the extended transition function is defined by:

$$\begin{aligned}\hat{\delta} &\in Q \times \Sigma^* \rightarrow Q \\ \hat{\delta}(q, \epsilon) &= q \\ \hat{\delta}(q, xw) &= \hat{\delta}(\delta(q, x), w)\end{aligned}$$

where  $q \in Q, x \in \Sigma, w \in \Sigma^*$ .