

G52MAL Machines and Their Languages Lecture 3

Non-deterministic Finite Automata (NFA)

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Recap: Language of a DFA

The **language** $L(A)$ defined by a DFA A is the set of words **accepted** by the DFA. For a DFA

$$A = (Q, \Sigma, \delta, q_0, F)$$

the language is defined by

$$L(A) = \{ w \in \Sigma^* \mid \hat{\delta}(q_0, w) \in F \}$$

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Recap: Formal Definition of DFA

Formally, a **Deterministic Finite Automaton** or **DFA** is defined by a 5-tuple

$$(Q, \Sigma, \delta, q_0, F)$$

where

- Q : **Finite** set of States
- Σ : Alphabet (finite set of symbols)
- $\delta \in Q \times \Sigma \rightarrow Q$: Transition Function
- $q_0 \in Q$: Initial or Start State
- $F \subseteq Q$: Accepting (or Final) States

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Recap: Extended Transition Function

The **Extended Transition Function** is defined on a state and a **word** (string of symbols) instead of on a single symbol.

For a DFA $A = (Q, \Sigma, \delta, q_0, F)$, the extended transition function is defined by:

$$\begin{aligned}\hat{\delta} &\in Q \times \Sigma^* \rightarrow Q \\ \hat{\delta}(q, \epsilon) &= q \\ \hat{\delta}(q, xw) &= \hat{\delta}(\delta(q, x), w)\end{aligned}$$

where $q \in Q, x \in \Sigma, w \in \Sigma^*$.

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