G52MAL Machines and their Languages Lecture 4: Nondeterministic Finite Automata (NFAs)

Thorsten Altenkirch based on slides by Neil Sculthorpe

Room A10 School of Computer Science University of Nottingham United Kingdom txa@cs.nott.ac.uk

8th February 2012

- DFA D = ($Q, \Sigma, \delta, q_0, F$), where:
 - Q: Finite set of states
 - Σ : Alphabet
 - $\delta \in (Q \times \Sigma \rightarrow Q)$: Transition function
 - $q_0 \in Q$: Initial state
 - $F \subseteq Q$: Set of accepting states
- Extended Transition Function:

$$egin{aligned} &\hat{\delta} \in (\mathcal{Q} imes \Sigma^* o \mathcal{Q}) \ &\hat{\delta} (q, arepsilon) &= q \ &\hat{\delta} (q, xw) &= \hat{\delta} (\delta (q, x), w) \end{aligned}$$

• Language of a DFA:

$$L(D) = \{ w \mid \hat{\delta}(q_0, w) \in F \}$$

(Informal) NFA Rules

Informal Rules

NFAs are mostly the same as DFAs, except:

- An NFA can "make choices".
- A word is accepted if choices can be made such that the machine ends in a final state.

Another way of looking at this is that:

- An NFA can be in multiple states at the same time.
- A word is accepted if any of the ending states are final.



Formal Definition of an NFA

A NFA N is a 5-tuple, $N = (Q, \Sigma, \delta, S, F)$, where:

- Q is a finite set of states
- Σ is an alphabet
- $\delta \in (Q \times \Sigma \rightarrow \mathcal{P}(Q))$ is a transition function
- $S \subseteq Q$ is a set of initial states
- $F \subseteq Q$ is a set of accepting (or final) states

The Extended Transition Function for NFAs

$$\begin{split} \hat{\delta} &\in (\mathcal{P}(Q) \times \Sigma^* \to \mathcal{P}(Q)) \\ \hat{\delta} & (P, \varepsilon) &= P \\ \hat{\delta} & (P, xw) &= \hat{\delta} & (\bigcup \{ \delta & (q, x) \mid q \in P \}, w) \end{split}$$

where

 $P \subseteq Q$ $x \in \Sigma$ $w \in \Sigma^*$

The Language of an NFA

Given an NFA $N = (Q, \Sigma, \delta, S, F)$, the language of N is defined as:

$$L(N) = \{ w \mid (\hat{\delta}(S, w) \cap F) \neq \emptyset \}$$

Recommended Reading

- Introduction to Automata Theory, Languages, and Computation (3rd edition), pages 55–59
- G52MAL Lecture Notes, pages 9–11